#### **How to Grow a TREE from CBASS**

### Interactive Binary Analysis for Security Professionals

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#### **Outline**

- Background
- Interactive Binary Analysis with TREE and CBASS
- Demonstrations
- Conclusions

# **Interactive Binary Analysis**

 Automated binary analyses useful for certain tasks (e.g., finding crashes)

Many binary analyses can't be automated

 Expert experience and heuristics are still key to binary analyses

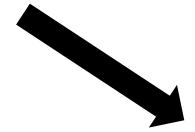
### **Benefits of Interactive Binary Analysis**

Applicable to many security problems

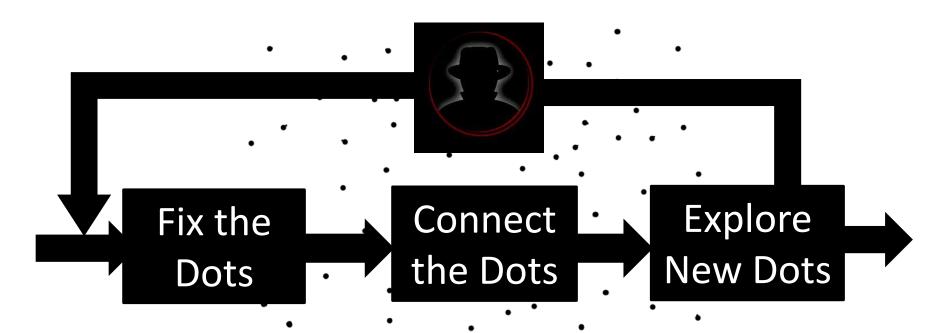
- Our tools increase productivity in:
  - Finding vulnerabilities
  - Analyzing root causes
  - Exploitability and risk assessment

#### **Interactive Analysis Like Connecting Dots**

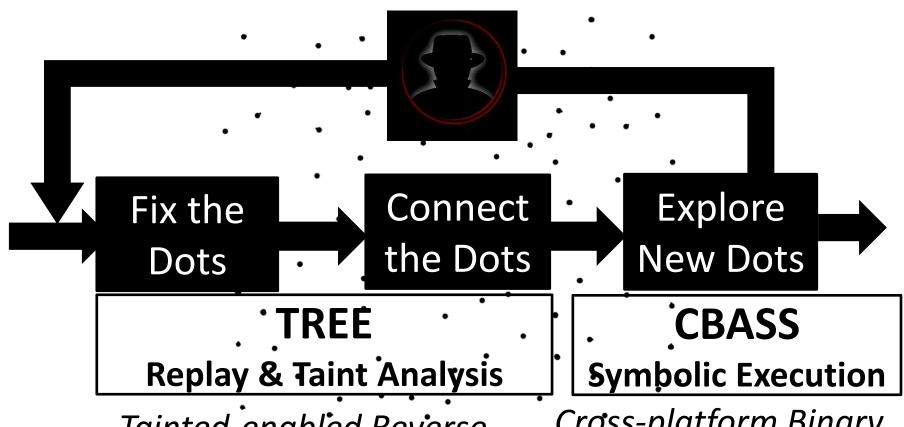




### Our Tools are Designed to Help



#### What Do Our Tools Do?



Tainted-enabled Reverse Engineering Environment Cross-platform Binary Automated Symbolic execution System

# Gaps between Research and Interactive Binary Analysis

- Existing research does not support interactive binary analysis
  - No practical tools
  - No uniform trace collection tools
  - No unified Instruction SetArchitecture(ISA) -independentanalysis tools

# Bringing Proven Research Techniques to Interactive Binary Analysis

- Our tools use dynamic, trace-based, offline analysis approach
  - Interactive binary analysis [1]
  - Dynamic taint analysis ([2][3][4])
  - Symbolic execution/ SMT solver ([2][5])
  - Trace replay ([6])

# **Making It Practical**

- TREE integrates with IDA Pro now and other mainstream binary analysis environments (later)
- TREE leverages debugging infrastructure to support tracing on multiple platforms
- CBASS uses Intermediate Representation (REIL [6][7])-based approach to support ISA-independent analysis

# CBASS Supports Both Automated & Interactive Analysis

**TREE** 

**Interactive Analysis** 

**Automated Fuzzer** 

**Automated Analysis** 

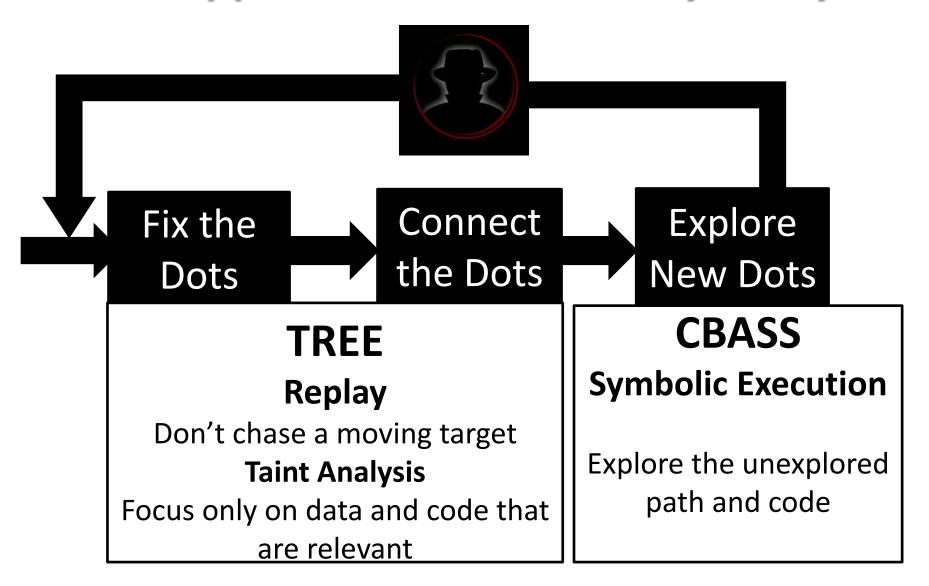
**CBASS** 

**IR-based Symbolic** 

**Execution Engine** 

TREE fills gaps for interactive analysis

### **Tools Support Interactive Binary Analyses**



# Illustrative Dots in Vulnerability Analysis: A Running Example

```
//INPUT
ReadFile(hFile, sBigBuf, 16, &dwBytesRead, NULL);
                                                 //Vulnerable Function
                                                 void StackOVflow(char *sBig,int num)
//INPUT TRANSFORMATIONS
                                                      char sBuf[8] = \{0\};
//PATH CONDITIONS
                                                      for(int i=0;i<num;i++)
if(sBigBuf[0]=='b') iCount++;
                                                      //Overflow when num>8
if(sBigBuf[1]=='a') iCount++;
if(sBiqBuf[2]=='d') iCount++;
                                                           sBuf[i] = sBig[i];
if(sBigBuf[3]=='!') iCount++;
if(iCount==4) // bad!
         StackOVflow(sBigBuf,dwBytesRead)
                                                      return;
else // Good
         printf("Good!");
```

### **Our Tools Support**

Fixing the Dots (TREE)

#### Fix the Dots

Reverse engineers don't like moving dots

- Why do the dots move?
  - Concurrency (multi-thread/multi-core)brings non-deterministic behavior
  - ASLR guarantees nothing will be the same

#### Fix the Dots

- How does TREE work?
  - Generates the trace at runtime
  - Replays it offline
- TREE trace
  - Captures program state = {Instruction, Thread,Register, Memory}
  - Fully automated generation
- TREE can collect traces from multiple platforms
  - Windows/Linux/Mac OS User/Kernel and real devices (Android/ARM, Cisco routers/MIPS, PowePC)

# TREE Taint-based Replay vs. Debug-based Replay

- Debug-replay lets you connect the dots
  - Single step, stop at function boundary, Breakpoint

- TREE replay connects dots for you
  - Deterministic replay with taint-point break

# **Our Tools Support**

**Connecting the Dots (TREE)** 

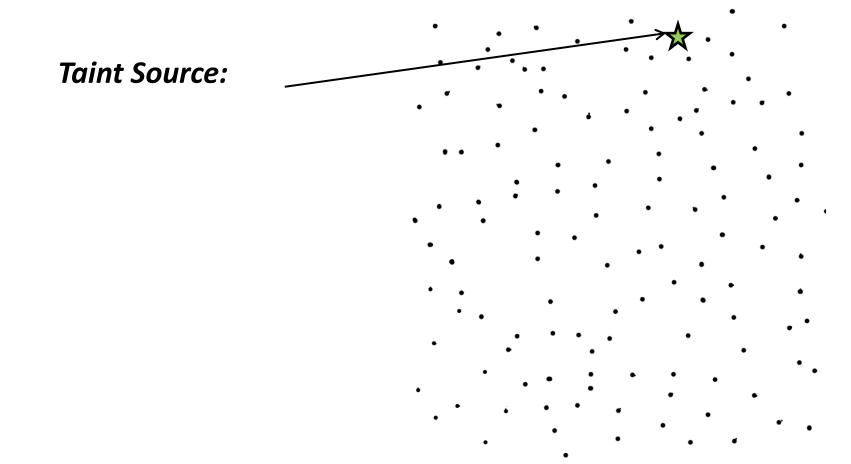
### **Connecting Dots is Hard**

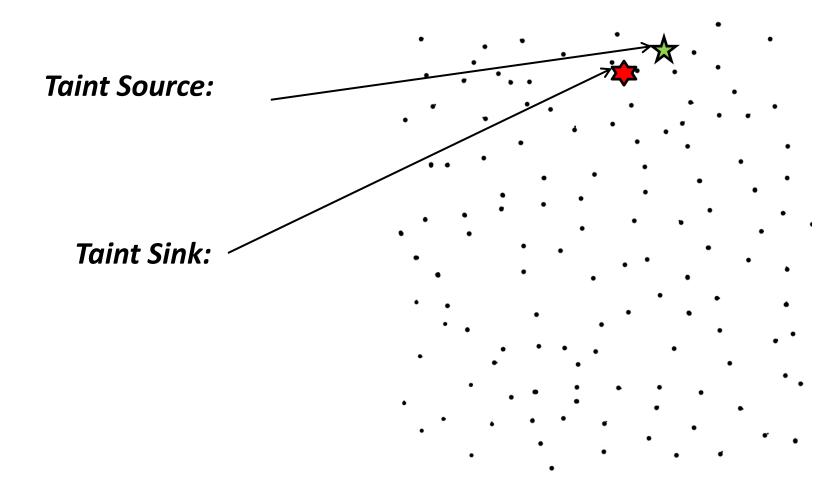
- Basic elements complex in real programs
  - Code size can be thousands (++) of lines
  - Inputs can come from many places
  - Transformations can be lengthy
  - Paths grow exponentially
- Basic elements likely separated by millions of instructions, spatially and temporally
- Multiple protections built in

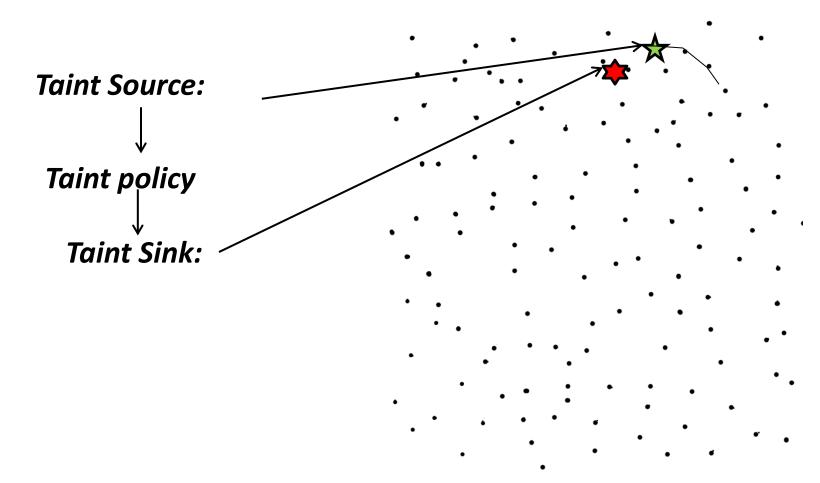
# **Techniques Help Connect the Dots**

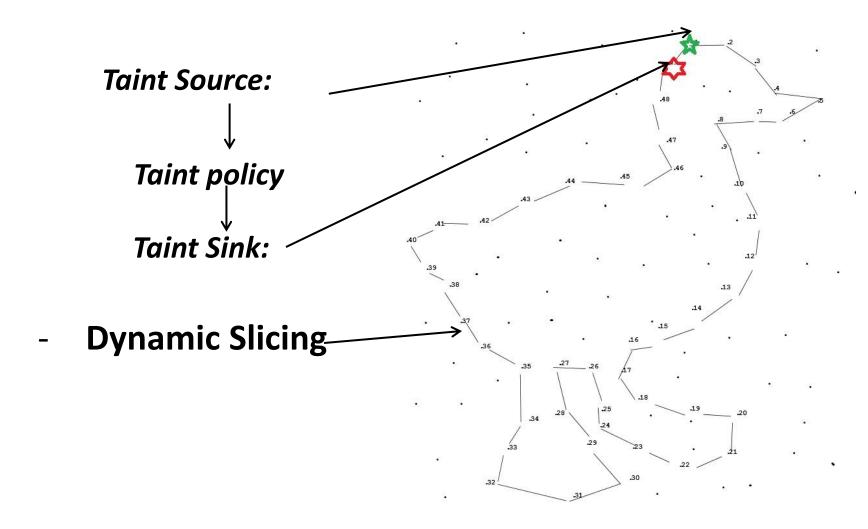
- Dynamic Taint Analysis
  - Basic Definitions
    - Taint source
    - Taint Sink:
    - o Taint Policy:

- Taint-based Dynamic Slicing
  - Taint focused on data
  - Slicing focused on relevant instructions and sequences





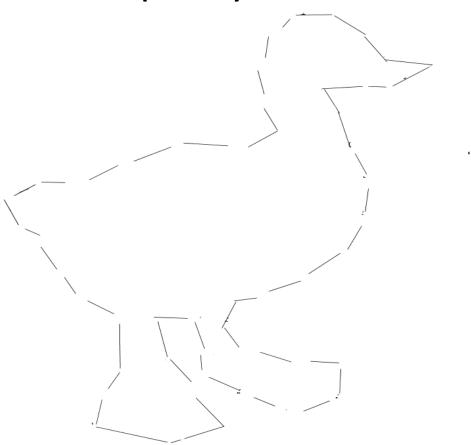




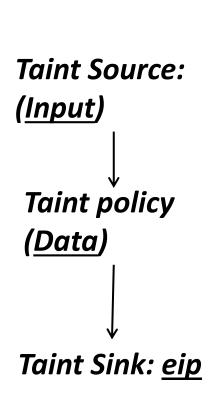
#### Find the Dots and Slice that Matter

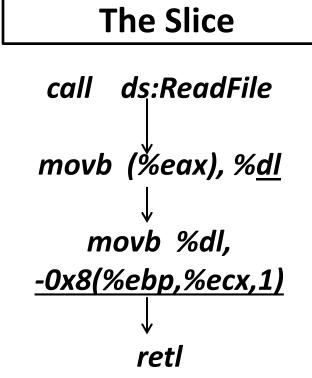
In practice, most dots don't matter – eliminate them quickly to focus on what

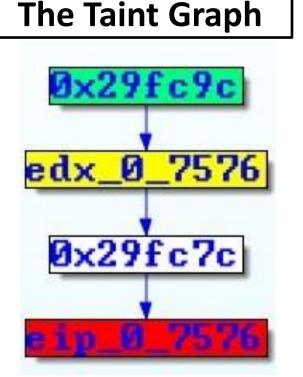
matters



### **Connecting Dots in Running Example**







#### What You Connect is What You Get

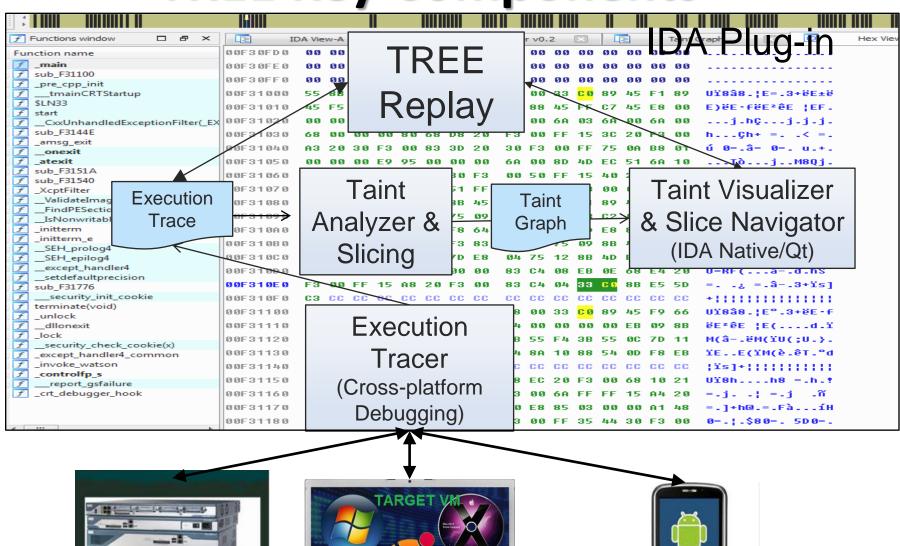
- Dots can be connected in different ways
  - Data dependency
  - Address dependency
  - Branch conditions
  - Loop counter

Connect dots in different taint policies



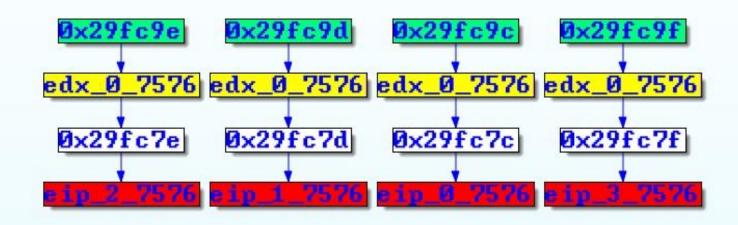
# TAINT-ENABLED

### **TREE Key Components**



CIORCID

Taint Graph

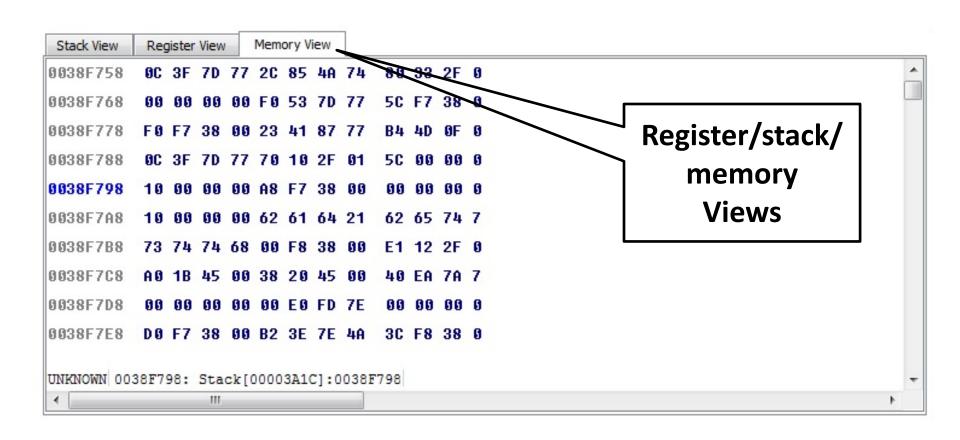


Taint Table

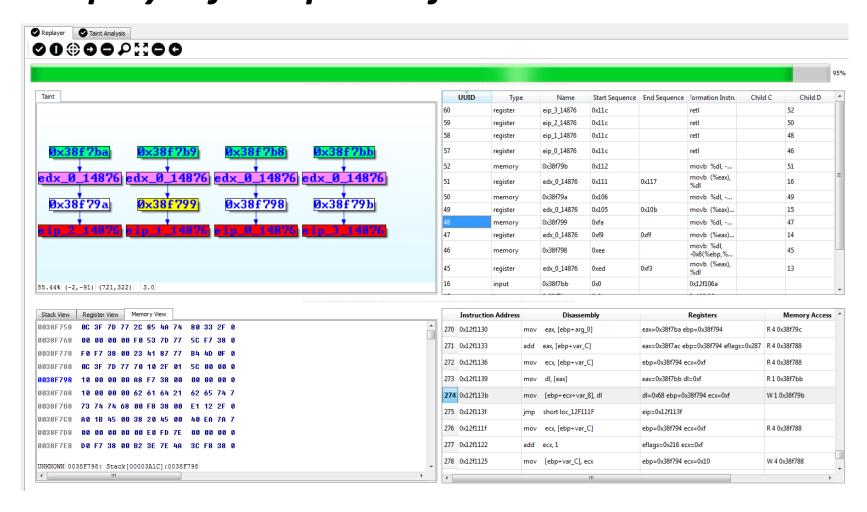
UUID	Туре	Name	Start Sequence	End Sequence	formation Instru	Child C	Child D	^
60	register	eip_3_14876	0x11c		retl		52	
59	register	eip_2_14876	0x11c		retl		50	
58	register	eip_1_14876	0x11c		retl		48	
57	register	eip_0_14876	0x11c		retl		46	
52	memory	0x38f79b	0x112		movb %dl,		51	
51	register	edx_0_14876	0x111	0x117	movb (%eax), %dl		16	Ш
50	memory	0x38f79a	0x106		movb %dl,		49	
49	register	edx_0_14876	0x105	0x10b	movb (%eax)		15	
48	memory	0x38f799	0xfa		movb %dl,		47	
47	register	edx_0_14876	0xf9	0xff	movb (%eax)		14	
46	memory	0x38f798	0xee		movb %dl, -0x8(%ebp,%		45	
45	register	edx_0_14876	0xed	0xf3	movb (%eax), %dl		13	
16	input	0x38f7bb	0x0		0x12f106a			
			12.2					

Execution
Trace
Table

	Instruction Address	Disassembly	Registers	Memory Access
270	0x12f1130	mov eax, [ebp+arg_0]	eax=0x38f7ba ebp=0x38f794	R 4 0x38f79c
271	0x12f1133	add eax, [ebp+var_C]	eax=0x38f7ac ebp=0x38f794 eflags=0x287	R 4 0x38f788
272	0x12f1136	mov ecx, [ebp+var_C]	ebp=0x38f794 ecx=0xf	R 4 0x38f788
273	0x12f1139	mov dl, [eax]	eax=0x38f7bb dl=0xf	R 1 0x38f7bb
274	0x12f113b	mov [ebp+ecx+var_8], dl	dl=0x68 ebp=0x38f794 ecx=0xf	W 1 0x38f79b
275	0x12f113f	jmp short loc_12F111F	eip=0x12f113f	
276	0x12f111f	mov ecx, [ebp+var_C]	ebp=0x38f794 ecx=0xf	R 4 0x38f788
277	0x12f1122	add ecx, 1	eflags=0x216 ecx=0xf	
278	0x12f1125	mov [ebp+var_C], ecx	ebp=0x38f794 ecx=0x10	W 4 0x38f788
4		111		



# TREE: The Front-end of Our Interactive Analysis System Replay is focal point of user interaction



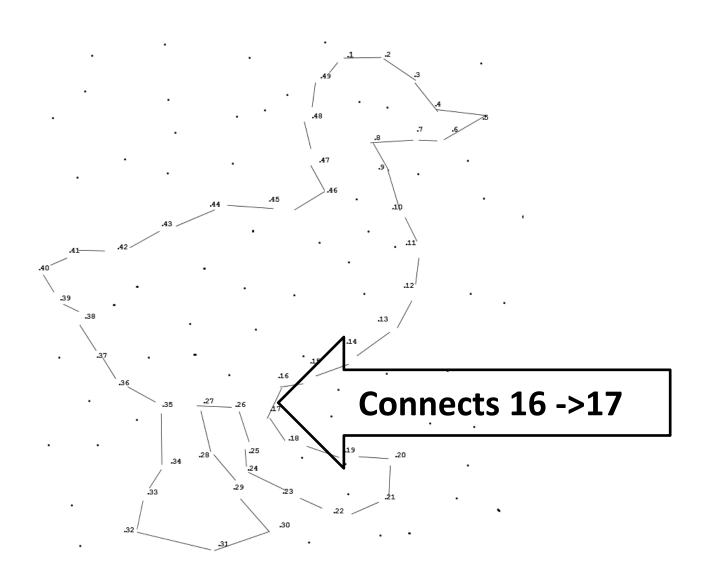
#### **Tree Demo**

# Using TREE to Analyze a Crash

## **Our Tools Support**

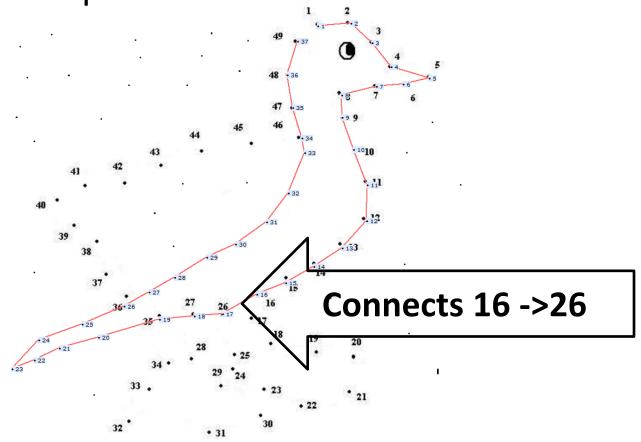
**Exploring New Dots** 

# A Key Branch Point for a Duck



# The Path for a ...

 Reverse engineers don't just connect dots; they want to explore new dots:



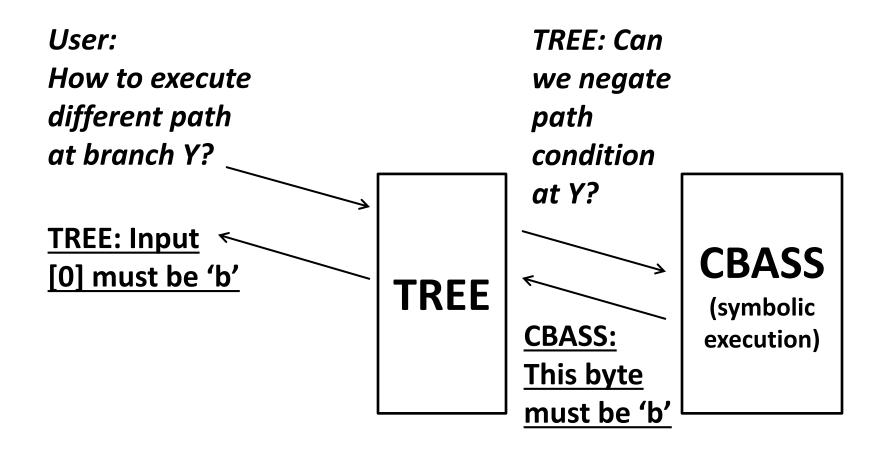
# **Explore New Dots**

 How do you force the program to take a different path to lead to "bad!"?

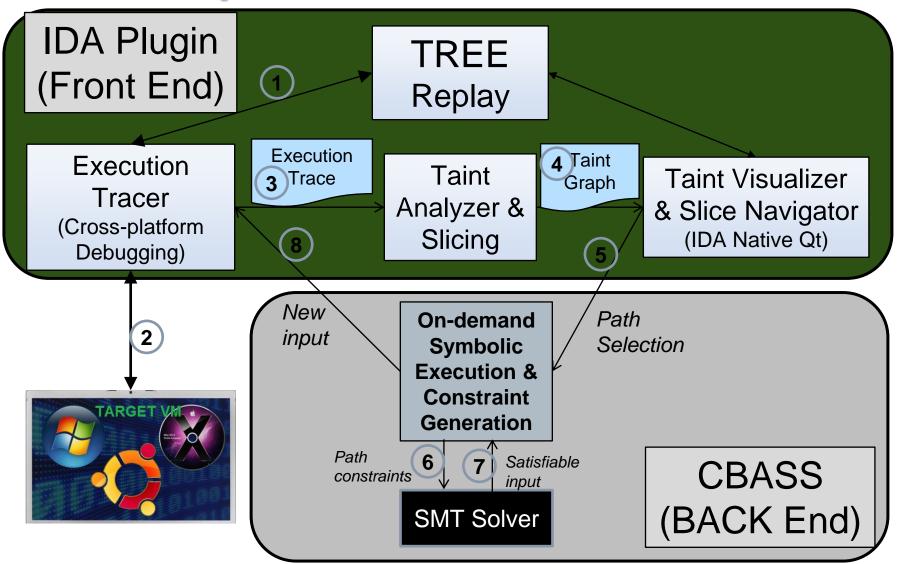
```
//INPUT
ReadFile(hFile, sBigBuf, 16, &dwBytesRead, NULL);
//PATH CONDITION
if(sBigBuf[0]=='b') iCount++;
if(sBigBuf[1]=='a') iCount++;
if(sBigBuf[2]=='d') iCount++;
if(sBigBuf[3]=='!') iCount++;
if(iCount==4) // "bad!" path
         StackOVflow(sBigBuf,dwBytesRead)?
Else // "Good" path
         printf("Good!");
```

# **Explore New Dots**

 User wants execution to take different path at a branch point Y – what input will make that happen?



# **Explore New Dots Demo**



# Task 1: Force the Program to Take "bad!" Path

#### //INPUT

ReadFile(hFile, sBigBuf, 16, &dwBytesRead, NULL);

#### //INPUT TRANSFORMATION

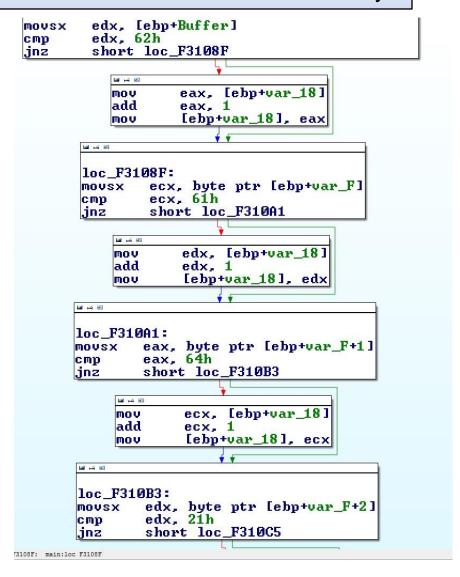
. . . . . .

#### //PATH CONDITION

```
if(sBigBuf[0]=='b') iCount++;
if(sBigBuf[1]=='a') iCount++;
if(sBigBuf[2]=='d') iCount++;
if(sBigBuf[3]=='!') iCount++;
if(iCount==4) // "bad!" path
//Vulnerable Function
```

StackOVflow(sBigBuf,dwBytesRead)
else
printf("Good!");

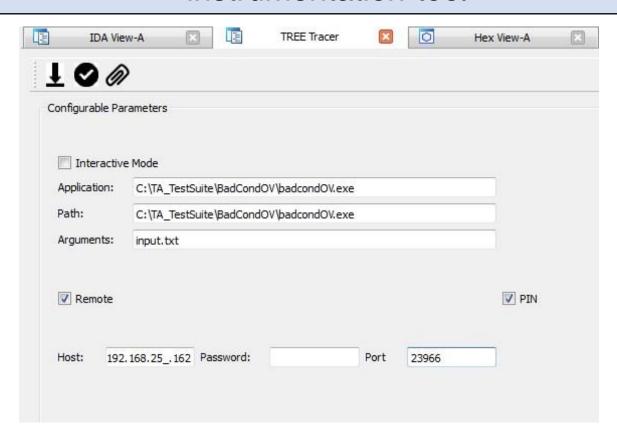
**Branch Conditions In Disassembly** 



### **1** TREE Pin Trace

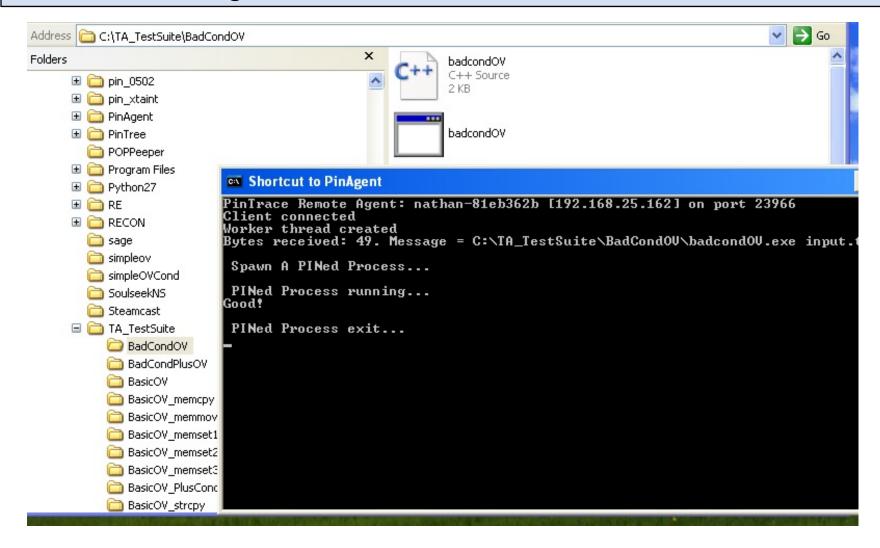
# PIN: A popular Dynamic Binary Instrumentation (DBI) Framework

http://software.intel.com/en-us/articles/pin-a-dynamic-binary-instrumentation-tool

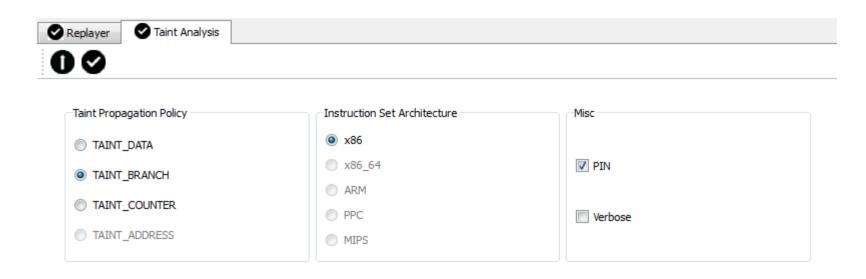


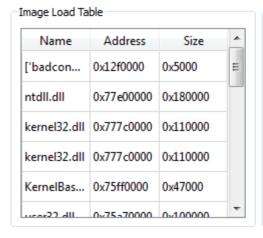
### 2 TREE Console: Trace Generation

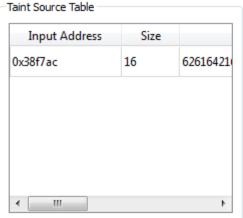
### PINAgent: Connects TREE with PIN tracer

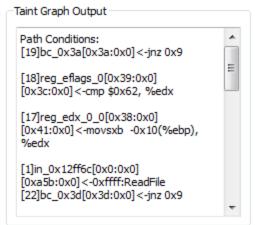


## 3 TREE: Taint Analysis Configuration

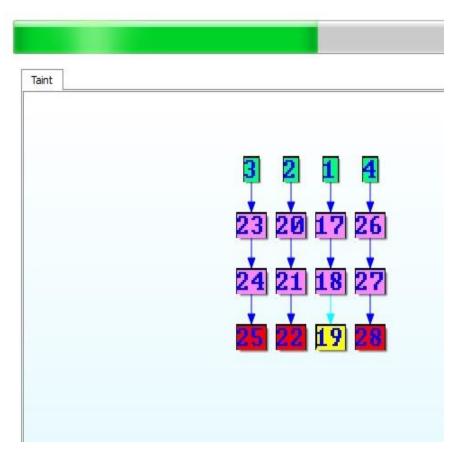






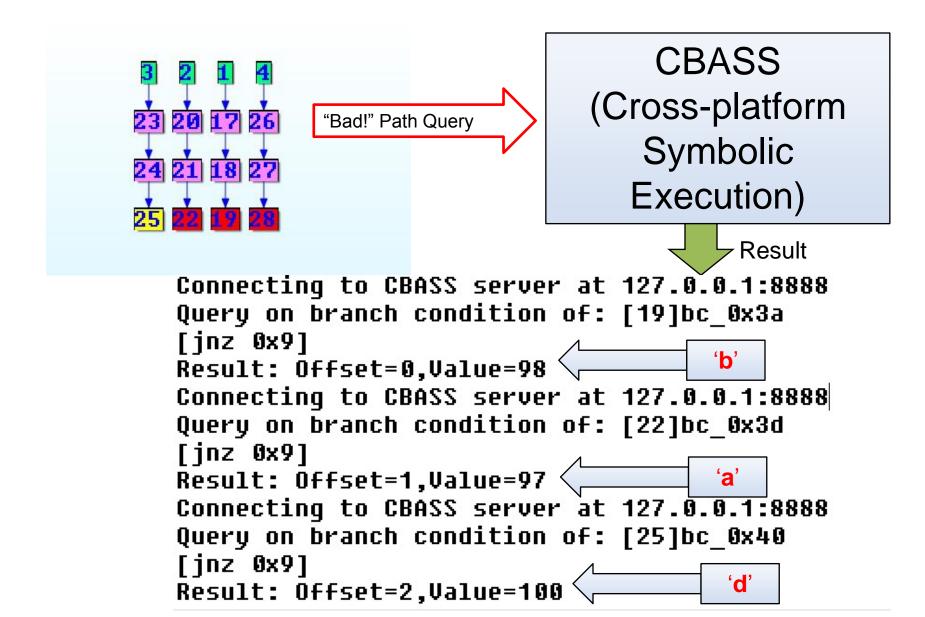


# 4 TREE: Branch Taint Graph



UÛID	Туре	Name	
1	input	0x12ff6c	
17	register	edx_0_0	
18	register	eflags 0	
19	branch	0x3a	
2	input	0x12ff6d	
20	register	ecx_0_0	
21	register	eflags_0	
22	branch	0x3d	
23	register	eax_0_0	
24	register	eflags_0	
25	branch	0x40	
26	register	edx_0_0	
27	register	eflags_0	
28	branch	0x43	
3	input	0x12ff6e	
input		0x12ff6f	

# ⑤ Negate Tainted Path Condition to Exercise a New ("Bad") Path



# On-demand Symbolic Execution (What Happens Behind the Scene)

```
(set-logic QF_AUFBV)

(declare-fun _IN_0x12ff6c_0x0_SEQ0 () (_ BitVec 8))
(declare-fun EXPR_0 () (_ BitVec 32))
(assert (= EXPR_0 (bvsub ((_ sign_extend 24) (bvxor_IN_0x12ff6c_0x0_SEQ0 (_ bv128 8))) (_ bv4294967168 32))))

(assert (= (ite (not (= (ite (not (= (bvand ((_ extract 63 0) (bvsub ((_ sign_extend 32) (bvand ((_ extract 31 0) EXPR_0) (_ bv4294967295 32))) (_ bv98 64))) (_ bv4294967295 64)) (_ bv0 64))) (_ bv1 32) (_ bv0 32)) (_ bv0 32))) (_ bv1 8) (_ bv0 8)) (_ check-sat)
(get-value (_IN_0x12ff6c_0x0_SEQ0))
```

```
RD t5, EMPTY , DWORD edx1]

58:481877 cmp edx, 98

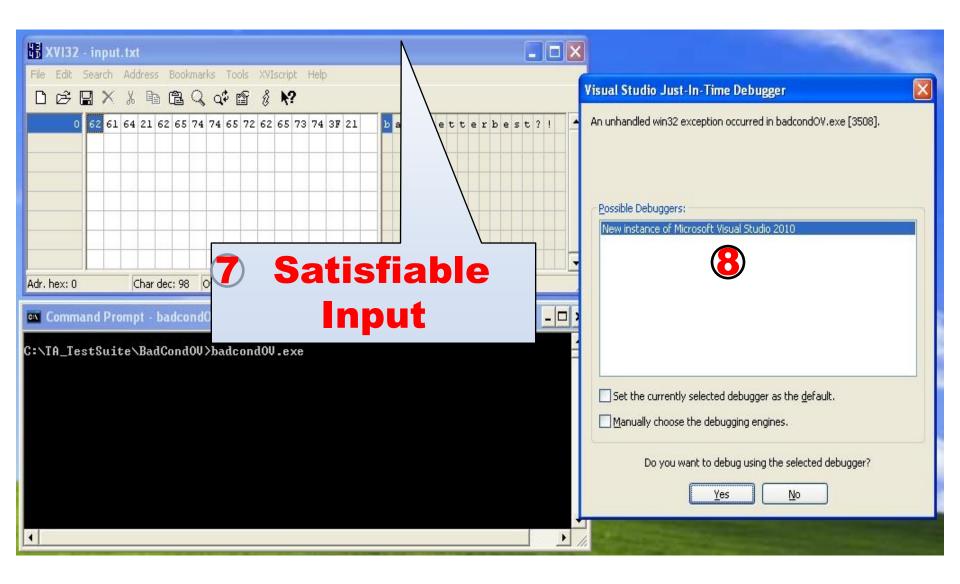
[48187708: and [DWORD edx, DWORD 2147483648, DWORD t8], 48187781: and [DWORD 98, DWORD 2147483648, DWORD t1], 48187782: sub [DWORD edx, DWORD 98, QWORD t2], 48187783: and [QWORD t2], QWORD 2147483648, DWORD t3], 48187784: bab [DWORD t3, DWORD -31, BYTE SF], 48187785: xor [DWORD t8, DWORD t1, DWORD t4], 48187786: xor [DWORD t0, DWORD t3, DWORD t5], 48187787: and [DWORD t4, DWORD t4], 48187786: xor [DWORD t0, DWORD t5], 48187787: and [DWORD t4, DWORD t5, DWORD t6], 48187788: bab [DWORD t2], 48187788: bab [DWORD t2], 48187788: bab [DWORD t2], 48187788: and [QWORD t2], 48187788: and [QWORD t2], 48187788: and [QWORD t2, QWORD 4294967295, DWORD t8], 4818778C: bisz [DWORD t8, EMPTY, BYTE ZF]]

59:48187a jnz loc_48188E

[481878 jnz loc_48188E

[
```

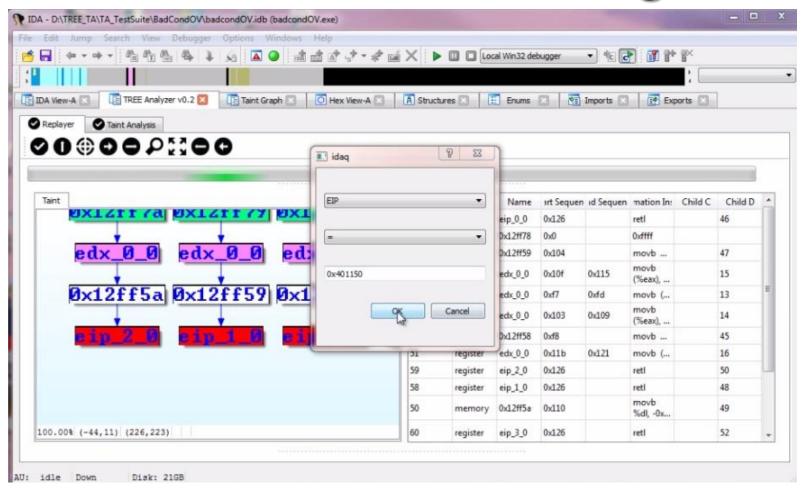
# TREE: Re-execute with Satisfiable Input



# Task 2: Own the Execution Assume Payload at 0x401150

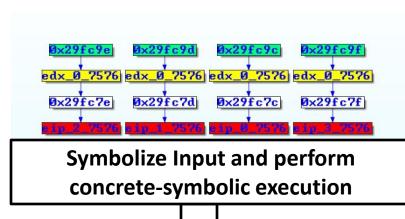
```
.text:00401145
                                align 10h
.text:00401150
                                 push
                                         ebp
.text:00401151
                                 mov
                                         ebp, esp
                                         1010h
.text:00401153
                                push
                                         offset aYouHaveBeenHac :
                                                                   "*** Yo
.text:00401158
                                 push
                                         offset aCbassCrossPlat :
.text:0040115D
                                push
.text:00401162
                                push
                                         ds:MessageBoxA
                                call
.text:00401164
.text:0040116A
                                push
                                         OFFFFFFFF
                                         ds exit
.text:0040116C
                                call
+av+•00401177
```

# **TREE Constraint Dialogue**



## Task 2:

# Own the Execution: From Crash to Exploit



### Symbolic eip =

(= expr\_0 (concat (bvand (bvor \_IN\_0x12ff6c\_0xd\_SEQ0 (\_ bv0 8)) (\_ bv255 8)) (bvand (bvor \_IN\_0x12ff6c\_0xc\_SEQ0 (\_ bv0 8)) (\_ bv255 8))))

#### **Query:**

get-value (\_IN\_0x12ff6c\_0xd\_SEQ0 \_IN\_0x12ff6c\_0xc\_SEQ0 \_IN\_0x12ff6c\_0xe\_SEQ0 \_IN\_0x12ff6c\_0xf\_SEQ0)

#### Sat:

(\_IN\_0x12ff6c\_0xd\_SEQ0 #x11 \_IN\_0x12ff6c\_0xc\_SEQ0 #x50 \_IN\_0x12ff6c\_0xe\_SEQ0 #x40 \_IN\_0x12ff6c\_0xf\_SEQ0 #x00

> SMT Solver

# TREE/CBASS Demo

# Using CBASS/TREE to Explore Bad Paths and Refine Exploits

### **Real World Case Studies**

Target Vulnerability	Vulnerability Name	Target Application Mode	Target OS
CVE-2005-4560	Windows WMF	User Mode	Windows
CVE-2207-0038	ANI Vulnerability	User Mode	Windows
OSVDB-2939	AudioCoder Vulnerability	User Mode	Windows
CVE-2011-1985	Win32k Kernel Null Pointer De- reference	Kernel Mode	Windows
CVE-2004-0557	Sound eXchange (SoX) WAV Multiple Buffer Overflow	User Mode	Linux
Compression/ Decompression	Zip on Android	User Mode	Real Device Trace Generation (In Progress)

# Highlights from Real World Case Study: Windows WMF Vulnerability (CVE-2005-4560)

- WMF SETABORTPROC Escape Vulnerability
  - http://www.cve.mitre.org/cgi-bin/cvename.cgi?name=CVE-2005-4560
  - The Windows Graphical Device Interface library (GDI32.DLL) in Microsoft Windows allows remote attackers to execute arbitrary code via a Windows Metafile (WMF) format image with a crafted SETABORTPROC GDI Escape function call, related to the Windows Picture and Fax Viewer (SHIMGVW.DLL).

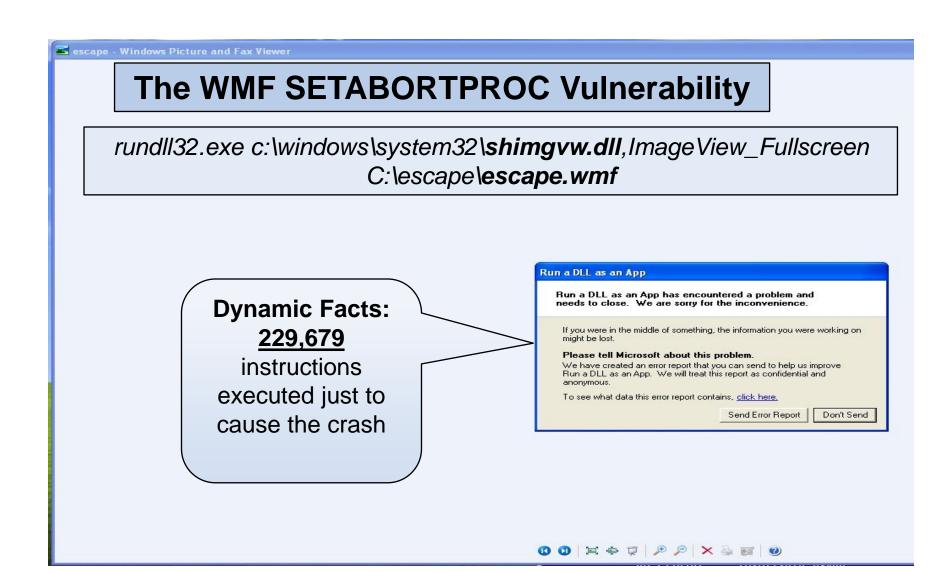
### **WMF Format**

- [MS-WMF]: Windows Metafile Format
  - http://msdn.microsoft.com/en-us/library/cc250370.aspx
- A Simplified One:
  - http://wvware.sourceforge.net/caolan/ora-wmf.html
- Overall WMF File Structure:

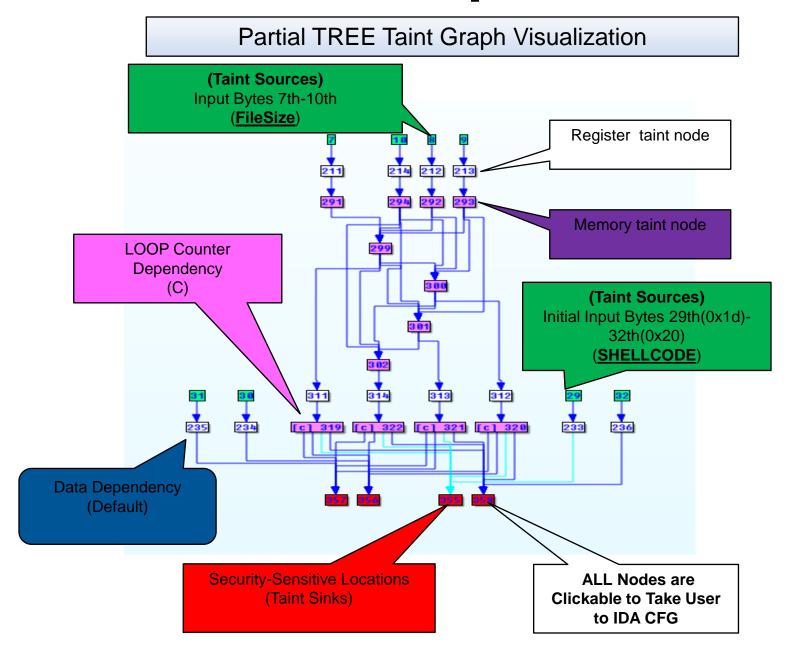
Meta	Meta	Meta	Meta	
Header	Record 1	Record 2	Record 3	

- One type of record is "escape" record
- SETABORTPROC escape allow an application to register a hook function to handle spooler errors

### **WMF Crash**



### **WMF Taint Graph**



## WMF File: The Fields & The Vulnerability

Key Structures:

```
typedef struct StandardMetaRecord
typedef struct WindowsMetaHeader
WORD FileType; /* Type of metafile (0=memory, 1=disk) */
                                                                 DWORD Size;
WORD HeaderSize: /* Size of header in WORDS (always 9) */
                                                                 /* Total size of the record in WORDs */
WORD Version: /* Version of Microsoft Windows used */
                                                                 WORD Function:
DWORD FileSize; /* Total size of the metafile in WORDs */
                                                                 /* Function number (defined in WINDOWS.H)
WORD NumOfObjects; /* Number of objects in the file */
DWORD MaxRecordSize; /* The size of largest record in WORDs */
                                                                 WORD Parameters[];
WORD NumOfParams; /* Not Used (always 0) */
                                                                 /* Parameter values passed to function */
                                                               } WMFRECORD:
} WMFHEAD;
           XVI32 - escape_wmf.NODISPLAY
            File Edit Search Address Bookmarks Tools XVIscript Help
            0 01 00 09 00 00 03 22 00 00 00 63 79 61 6E 69 64 2D 45 01 04 00 00 00 06 26 09 CC CC CC
                 00|00|00|80|03|00|00|00|00|00
Shellcode
                                                                                             SetAbortProc
                                                                             Escape
                           Char dec: 0
                                    Overwrite
           Adr. hex: 43
```

## WMF Slicing (1)

#### An Instruction Slice Traced Back from Crash Site to Input

Each node uniquely trace back to one execution event through its sequence number

0x77f330a3 <u>call eax</u> 2 ffd0 0x0 0x3812f Reg( EAX=0xa8b94 ESP=0xb4fb88 EIP=0x77f330a3 ) W 4 b4fb88

0x77c472e3 <u>rep movsd</u> 2 f3a5 0x0 0xb142 Reg( EDI=0xa8804 eflags=0x10216 ESI=0xa9f8c ECX=0xa ) R 4 a9f8c cc\_cc\_cc W 4 a8804

0x77f2e997 mov ecx, [ebp+arg\_8] 3 8b4d10 0x0 0xc5c3 Reg( EBP=0xb4fbf8 ECX=0x7c809a20 ) R 4 b4fc08 44\_0\_0

0x77f2e983 mov [ebp+arg\_8], eax\_ 3 894510 0x0 0xbd8c Reg( EAX=0x44 EBP=0xb4fbf8 ) W 4 b4fc08

0x77f2e97f add eax, eax 2 03c0 0x0 0xbd89 Reg( EAX=0x22 eflags=0x246 )

0x77f2e949 mov eax, [edi+6] 3 8b4706 0x0 0xbd7d Reg( EAX=0xa8920 EDI=0xa87e8 ) R 4 a87ee 22\_0\_0

0x77c472e3 <u>rep movsd</u> 2 f3a5 0x0 0xb13c Reg( EDI=0xa87ec eflags=0x10216 ESI=0xa9f74 ECX=0x10 ) R 4 a9f74 0\_3\_22\_0 W 4 a87ec

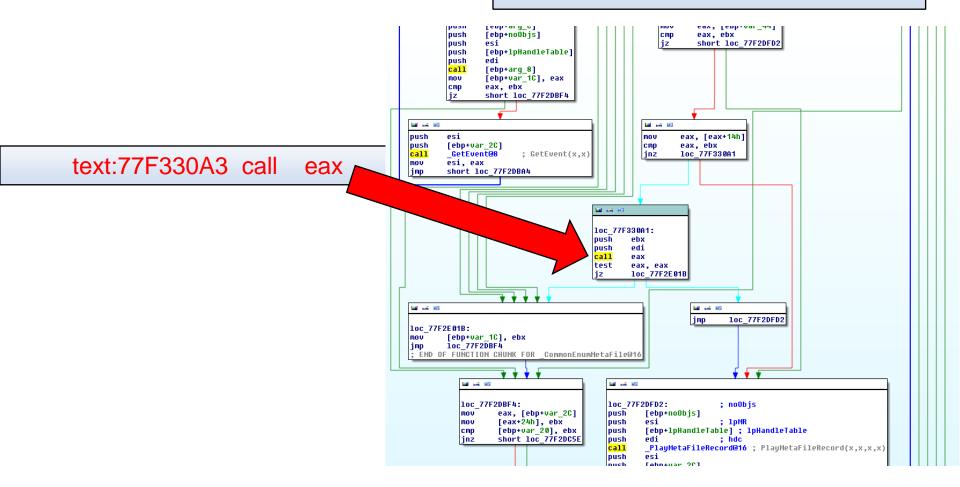
## WMF Slicing (2)

An Instruction Slice with Text Helps

Put Instruction In Its Context Helps
More

Module: gdi32.dll

Function: CommonEnumMetaFile



## WMF Slicing (3)

An Instruction Slice with Text Helps a Little

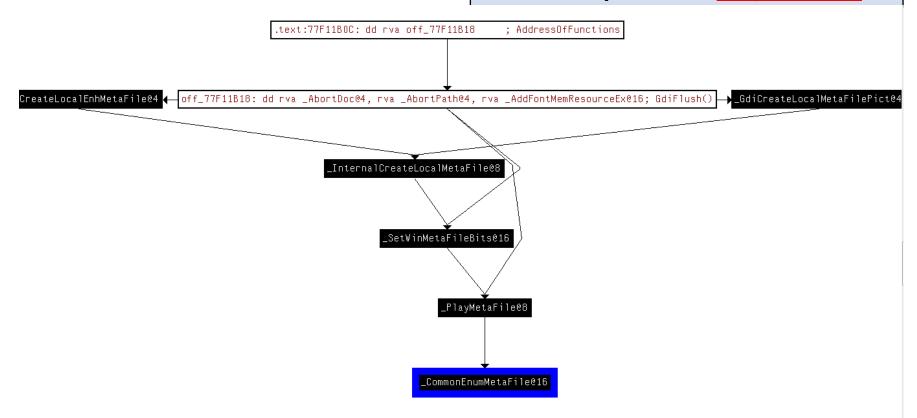
text:77F330A3 call eax

More Context Helps More

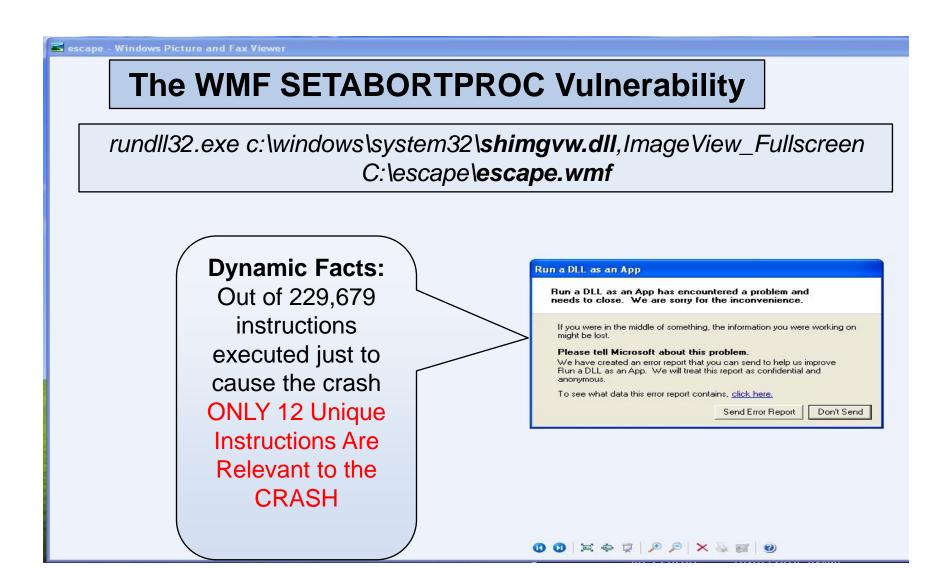
Module: gdi32.dll

Function: CommonEnumMetaFile

Call Graph: caller **PlayMetaFile** 



### **WMF** -- The Relevant Parts



# Conclusions

- Our tools support interactive binary analysis, with Replay, Dynamic Taint Analysis, and Symbolic Execution.
- TREE runs on top of IDA Pro and supports cross-platform trace collection, taint analysis and replay.
- CBASS (based on REIL) enables IR-based architecture-independent symbolic execution and can support both automated and interactive analysis.
- YOU drive the tools!

# Where You Can Get TREE

- TREE is open source at: http://code.google.com/p/tree-cbass/
  - First version of TREE (Taint Analysis) is released
  - Replay is in Progress
  - CBASS is Following
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