

dirtbox, a x86/Windows Emulator

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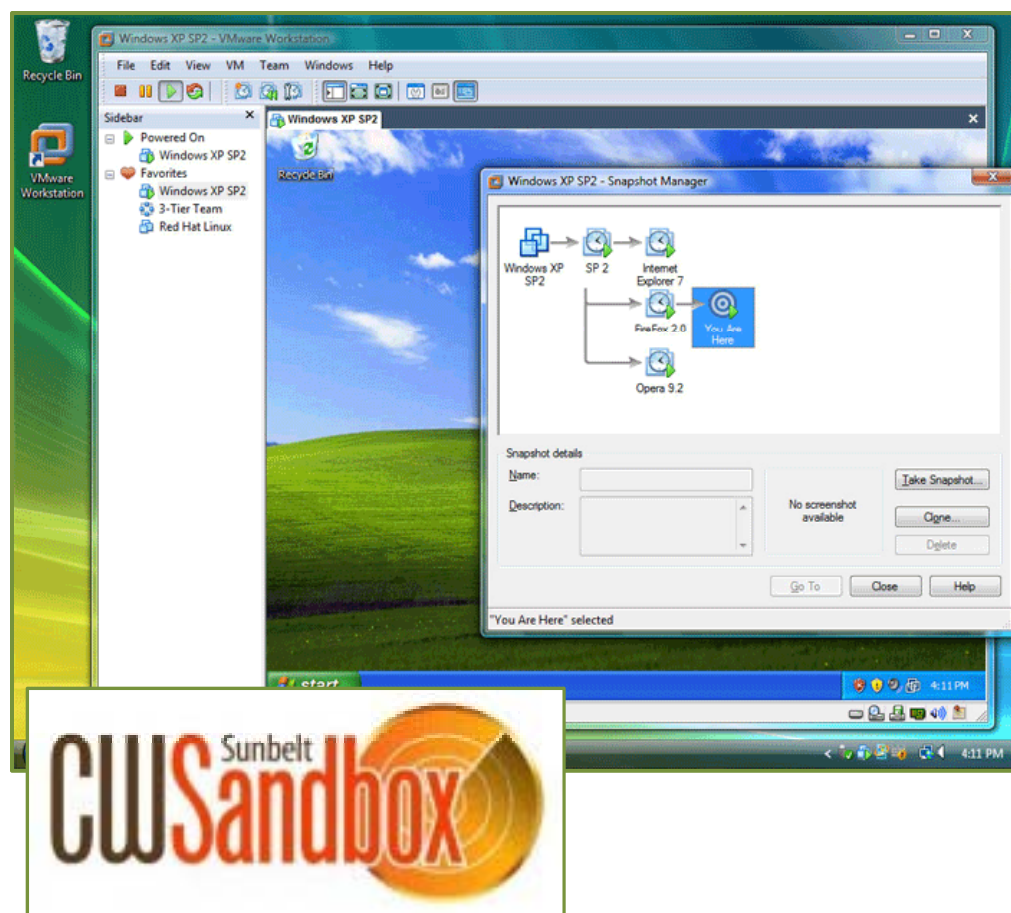
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BlackHat USA, 2010-07-29

Motivation & System Overview

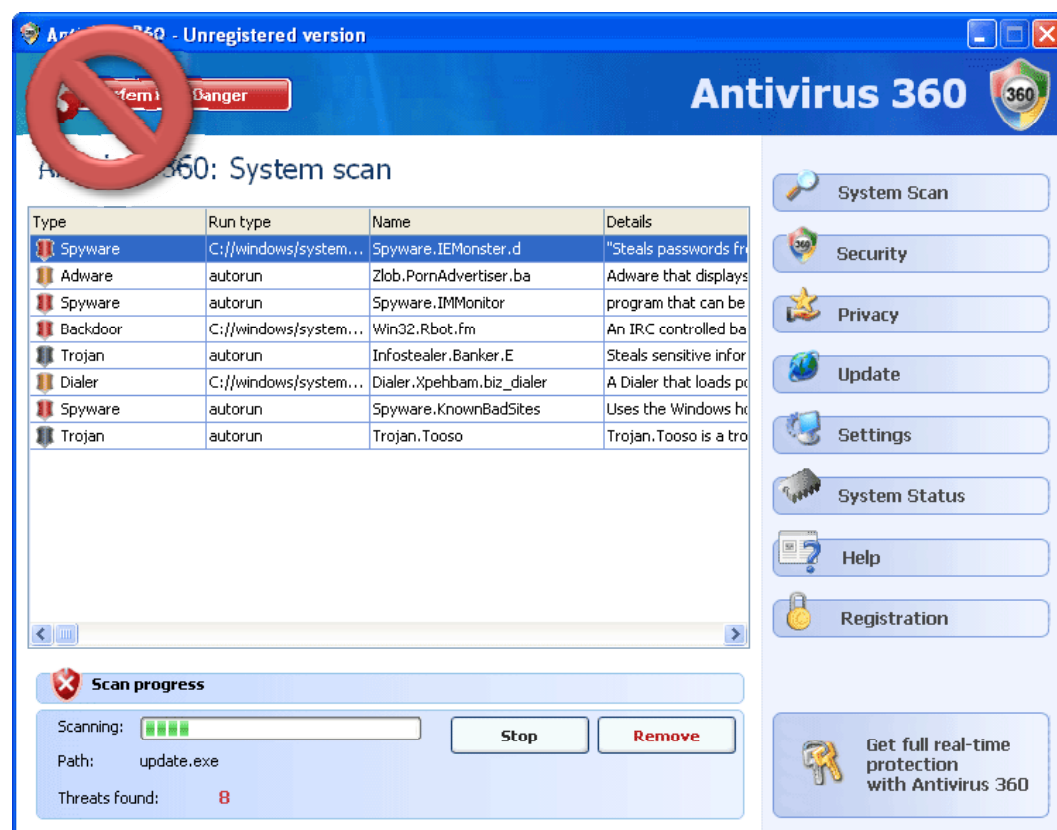
Why not just use CWSandbox, Anubis, Norman's, JoeBox, ...

Malware Analysis Sandbox Solutions



- VMWare „Rootkits“
 - CWSandbox
 - JoeBox
 - ThreatExpert
 - zBox
 - ...
- Norman Sandbox
- Anubis

Malware *Detection* Emulators (A/V)



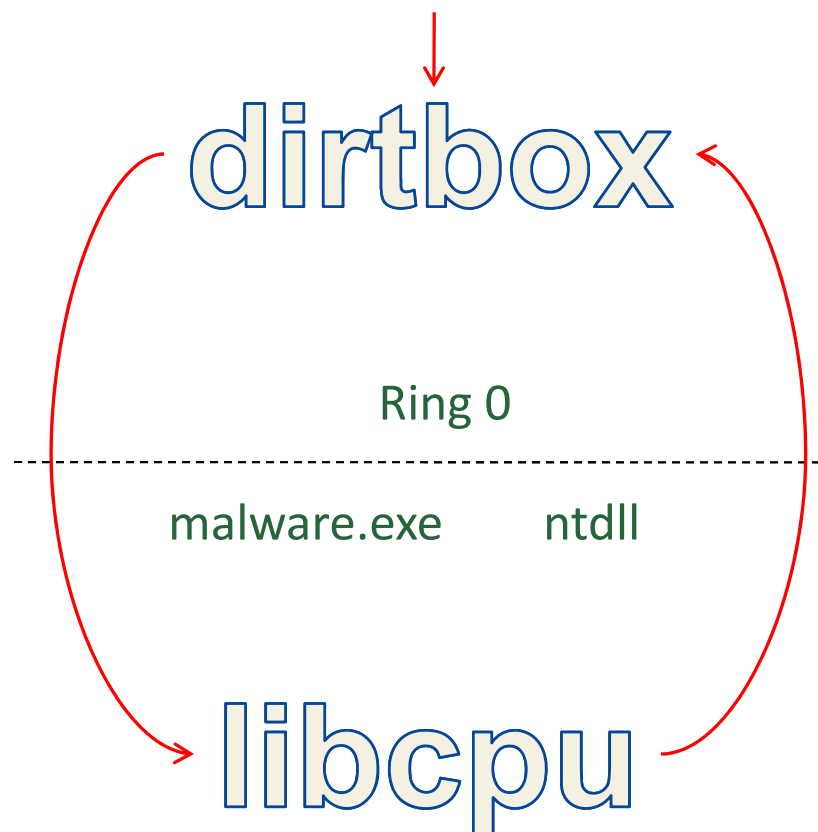
- Most *serious* A/V solutions have one
- API level emulation
- Often pure software emulators
- Detection by
 - Unimplented APIs
 - Heap Layout, SEH handling, ...
 - ...

- Functions containing `try {` in VS C++ share code
 - Epilogue is always the same
 - Uses sequence `push ecx / ret` to return to caller
 - The `ecx` register belongs to the called function by definition, so it is undefined upon API return
 - The `ecx` value can be predicted because it will point to the API's `ret`
- This breaks *a lot* of A/V emulators right away
 - There are some funny but trivially detected workarounds
 - Could be used for generic anti-emulation detection (use of undefined registers after SEH protected API calls)
- Relies on the fact that the API's bytecode is not emulated

System Overview or „A cat pooped into my sandbox and now I have a dirtbox!“



- System Call Layer
Emulation of Windows
- ntdll's native code is run inside virtual CPU
 - Other libraries wrap around kernel32 which wraps around ntdll
- Malware issuing system calls directly supported

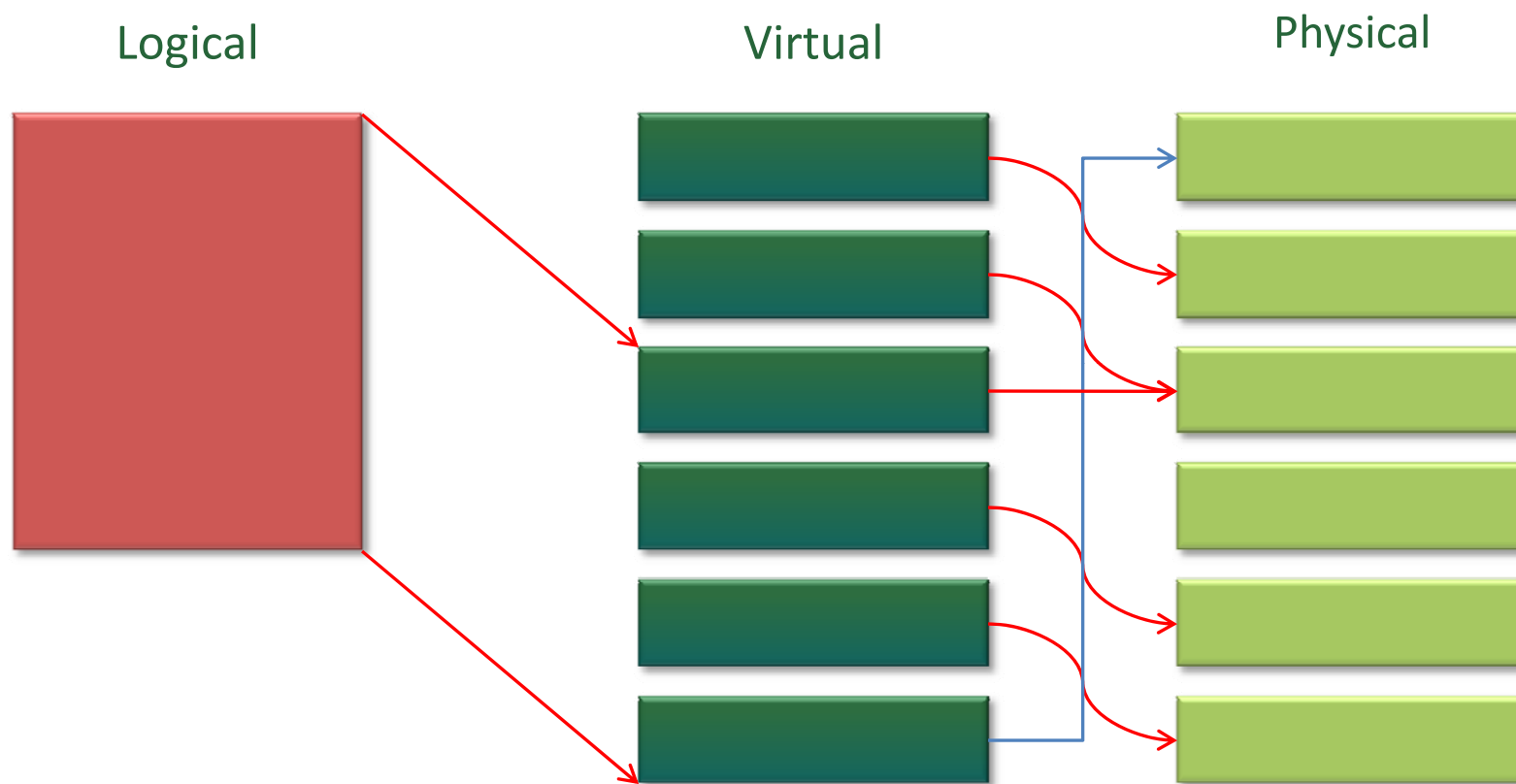


libcpu

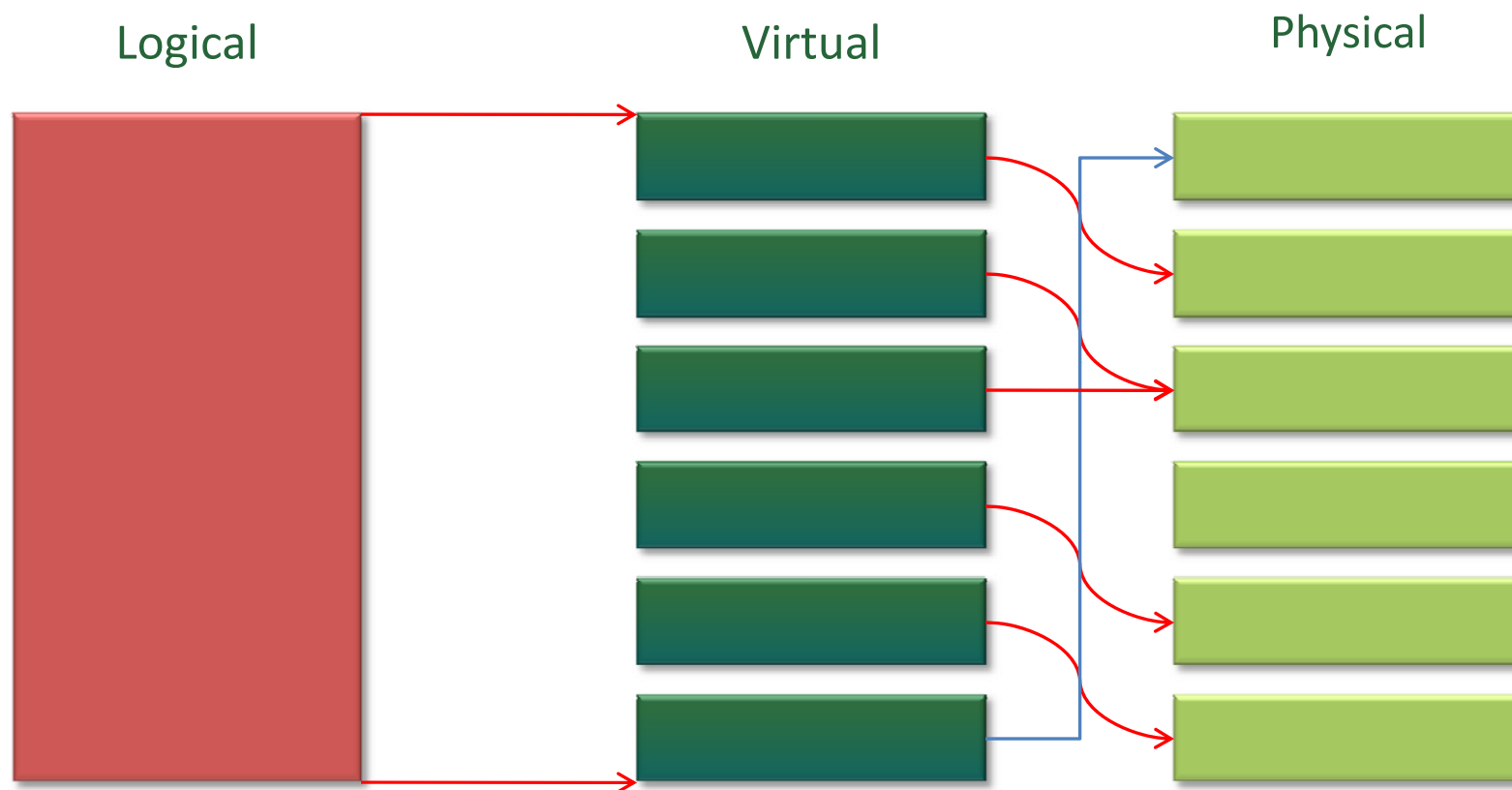
Custom x86 Basic Block Level Virtualization

- Software emulation of x86 bytecode is too slow
 - A lot of additional code, such as ntdll & kernel32
 - Existing Virtualization solutions are too powerful
 - Implementing their own MMU, support for privileged instructions
 - We want instruction level introspection
- Homebrew x86 virtualization based on LDT

x86 Memory Views



x86 Memory View on Current OS



- Global Descriptor Table
 - Allocated by Operating System
 - Shared among processes
 - Local Descriptor Table
 - Has to be allocated by the OS, too
 - SYS_modify_ldt
 - NtSetLdtEntries
 - Process specific, usually not present
- Define 2 GB guest „userland“ LDT segment

- Basic block level execution on host CPU
 - No instruction rewriting required (thanks to host MMU)
- Basic block is terminated by
 - Control flow modifying instruction
 - Privileged instructions
- Exception: Backward pointing jumps
 - Directly copy if points into same basic block
 - Enhanced loop execution speeds
- Currently no code cache, could cache disassembly results (length of basic block)

Self-Modifying Code



~PROT_WRITE

libcpu Demo



```
edi: 00000000
esi: 00000000
ebp: 00000000
esp: 00400000
ebx: 00000000
edx: 00000000
ecx: 00000000
eax: 00000000
eflags: 00000000
eip: 00400000
00400000 > mov ecx,0x1000000
00400005 > xor eax,eax
00400007 > inc eax
00400008 > dec ecx
00400009 > jnz 0x400007
0040000b <
edi: 00000000
esi: 00000000
ebp: 00000000
esp: 00400000
ebx: 00000000
edx: 00000000
ecx: 00000000
eax: 01000000
eflags: 00000246
eip: 0040000b
---
edi: 00000000
esi: 00000000
ebp: 00000000
esp: 00400000
ebx: 00000000
edx: 00000000
ecx: 00000000
eax: 01000000
eflags: 00000246
eip: 0040000b
Emulating 0040000b: int3
```



The word "libscizzle" is written in a large, bold, black, sans-serif font. It is positioned in the lower-left quadrant of the slide, above the subtitle. The background of the slide features a grey gradient with abstract, flowing lines and a series of white vertical bars of varying heights at the bottom, resembling a stylized city skyline or data visualization.

Or „libx86shellcodedetection“ if you prefer...

- Simple Approach: Brute-Force over byte buffer
 - If n valid instructions can be executed from there, assume we found valid shellcode
- Pre-filter buffers: Scan for „GetPC“ sequences
 1. Find GetPC opcode candidates: 89, a3, d9, e8
 - `mov r/m32, r32` or `mov rm/32, eax` → SEH based GetPC
 - `fstenv`
 - `call rel32`
 2. Check for valid memory operands or FS prefix
 - Require `fstenv` operand to be `esp` relative

Free Shellcode *Detector*:
<http://code.mwcollect.org/libscizzle>

Free Shellcode *Emulator*:
<http://libemu.carnivore.it/>

libscizzle Demo



```
edi: 00000000
esi: 00000000
ebp: 00000000
esp: 00400000
ebx: 00000000
edx: 00000000
ecx: 00000000
eax: 00000000
eflags: 00000000
eip: 00400000
00400000 > mov ecx,0x1000000
00400005 > xor eax,eax
00400007 > inc eax
00400008 > dec ecx
00400009 > jnz 0x400007
0040000b <
edi: 00000000
esi: 00000000
ebp: 00000000
esp: 00400000
ebx: 00000000
edx: 00000000
ecx: 00000000
eax: 01000000
eflags: 00000246
eip: 0040000b
---
edi: 00000000
esi: 00000000
ebp: 00000000
esp: 00400000
ebx: 00000000
edx: 00000000
ecx: 00000000
eax: 01000000
eflags: 00000246
eip: 0040000b
Emulating 0040000b: int3
```



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The background of the slide is a grayscale abstract image. It features a city skyline with various building heights in the foreground. In the background, there is a large, rounded dome-like structure, possibly representing a stadium or a large architectural feature. The overall aesthetic is modern and architectural.

dirtbox

Or „The System Call Implementor’s Sysiphus Tale“

- System Calls mostly undocumented
 - Wine, ReactOS, ...
- We get a lot of genuine environment for free!
- There is a fixed number of system calls but an unbound number of APIs (think third party DLLs)
- Some malware uses system calls directly anyway
- Less detectability by API side effects (because we run original bytecode)

- Process startup handled mostly by new process
 - Creating process allocates new process: NtCreateProcess
 - Creates „Section“ of new image & ntdll and maps into process, this requires kernel to parse section headers
 - Creates new Thread on Entry Point with APC in ntdll
 - ntdll!LdrInitializeThunk will relocate images if necessary, resolve imports recursively, invoke TLS and DLL startup routines and do magic (see demo).
- All we have to implement is NtCreateSection & NtMapViewOfSection for SEC_IMAGE → we only need to parse PE's section headers!

- A lot of A/V emulators naturally come with their own guest heap allocator implementations
 - Some even do not put heap headers before blocks
 - Let alone arena structures, ...
- The Windows heap is implemented in ntdll
 - Interfacing the kernel with NtVirtualAlloc & NtVirtualFree
 - All protections like heap cookies are present
- Fingerprinting other emulators:
 - Look at `malloc(0)-8`, look for proper block header
 - Or overflow until the heap cookie and free

- Generate **CONTEXT** record from current CPU state
- Jump to `ntdll!KiUserExceptionDispatcher`
- `ntdll` will do proper SEH handling for us
 - Lookup current top of SEH chain in TEB
 - Walk list, invoke exception handlers with correct flags
 - Checking for SafeSEH structures etc.
- Trivial detection for other emulators:
 - Link with SafeSEH header
 - Trigger exception with invalid handler registered
 - Check in `UnhandledExceptionHandler`

dirtbox Demo



```
src/win32/Process.cpp:0917 systemCall: Invoking system call #170: NtQuerySymbolicLinkObject
src/win32/Process.cpp:0592 sysc_NtQuerySymbolicLinkObject: <7c980280, 0>, 0 -> 'C:\WINDOWS\system32\'
src/win32/Process.cpp:0917 systemCall: Invoking system call #25: NtClose
src/win32/Process.cpp:0121 sysc_NtClose: Handle #c
DEBUG LDR: NEW PROCESS
DEBUG Image Path: C:\WINDOWS\system32\simple-time.exe (simple-time.exe)
DEBUG Current Directory: C:\WINDOWS\system32\
DEBUG Search Path: C:\WINDOWS\system32;C:\WINDOWS\system;C:\WINDOWS;.
src/win32/Process.cpp:0917 systemCall: Invoking system call #116: NtOpenFile
src/win32/Process.cpp:1042 resolveObjectAttributes: (18, 0, 9d84<'\\??\C:\WINDOWS\system32\'>, 42)
src/win32/Process.cpp:0374 sysc_NtOpenObject: Allocated handle 10 for '\\??\C:\WINDOWS\system32\'
src/win32/Process.cpp:0917 systemCall: Invoking system call #179: NtQueryVolumeInformationFile
src/win32/Process.cpp:0917 systemCall: Invoking system call #83: NtFreeVirtualMemory
src/win32/Process.cpp:0917 systemCall: Invoking system call #17: NtAllocateVirtualMemory
src/win32/Process.cpp:0045 sysc_NtAllocateVirtualMemory: ffffffff, 23000, -, 1000, 1000, 4
src/win32/Process.cpp:0917 systemCall: Invoking system call #84: NtFsControlFile
src/win32/Process.cpp:0446 sysc_NtFsControlFile: <10, 0, 0, 0, 99d8, 90028, (nil), 0, (nil), 0>
src/win32/Process.cpp:0917 systemCall: Invoking system call #139: NtQueryAttributesFile
src/win32/Process.cpp:1042 resolveObjectAttributes: (18, 0, 9db8<'\\??\C:\WINDOWS\system32\simple-time.exe.Local\'>, 40)
src/win32/Process.cpp:1055 resolveObjectAttributes: Object name could not be resolved: '\\??\C:\WINDOWS\system32\simple-time.exe.Local\'
src/win32/Process.cpp:0917 systemCall: Invoking system call #83: NtFreeVirtualMemory
DEBUG LDR: LdrLoadDll, loading kernel32.dll from
src/win32/Process.cpp:0917 systemCall: Invoking system call #125: NtOpenSection
src/win32/Process.cpp:1042 resolveObjectAttributes: (18, 8, 9a24<'kernel32.dll\'>, 40)
src/win32/Process.cpp:0374 sysc_NtOpenObject: Allocated handle 14 for 'kernel32.dll'
src/win32/Process.cpp:0917 systemCall: Invoking system call #108: NtMapViewOfSection
src/win32/Process.cpp:0322 sysc_NtMapViewOfSection: (14, ffffffff, 0, 0, 0, (nil), * 0x77445aec = 0, 1, 0, 4)
src/win32/SectionObject.cpp:0109 mapView: Successful PE loading: 7c800000, f6000
src/win32/Process.cpp:0917 systemCall: Invoking system call #25: NtClose
src/win32/Process.cpp:0121 sysc_NtClose: Handle #14
DEBUG LDR: ntdll.dll used by kernel32.dll
DEBUG LDR: Snapping imports for kernel32.dll from ntdll.dll
src/win32/Process.cpp:0917 systemCall: Invoking system call #137: NtProtectVirtualMemory
src/win32/Process.cpp:0238 sysc_NtProtectVirtualMemory: ffffffff, 7c801000, 624, 4 -> 3, 97bc
src/win32/Process.cpp:0917 systemCall: Invoking system call #137: NtProtectVirtualMemory
src/win32/Process.cpp:0238 sysc_NtProtectVirtualMemory: ffffffff, 7c801000, 1000, 20 -> 5, 20
src/win32/Process.cpp:0907 systemCall: Unsupported system call #78: NtFlushInstructionCache!
DEBUG LDR: LdrGetProcedureAddress by
```

Conclusion & Future Work

Let's use this for exploit development!

- No leaked registers in Ring 0 transition except for `eax`
 - Need to provide proper return codes, esp. error codes
 - `ntdll` just cares about $\geq 0xc0000000$; malware might look for specific error codes
- Side effects on buffers etc., especially in error cases
 - Fill out `IN OUT PDWORD Length` in case of error?
 - Roll back system calls performing multiple things?
- Tradeoff between detectability and performance

Future Work: Adding Tainting & SAT Checking



- Already did Proof-of-Concept based on STP
- Interleave static analysis into dynamic emulation
 - Look for interesting values (e.g. reads from network, date)
 - Do static forward data-flow analysis on usage
 - If used in conditional jumps, identify interesting values with a SAT Checker (there are better domain specific ways, but I'm lazy)
- Automatic reconstruction of network protocols (e.g. commands in IRC bots)
- Identify specific trigger based behaviour
- Identify Anti-Emulation behaviour

Questions? Thank You!

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